Updated 10/31/2022

Topic 2: Database design L15: ER & Relational modeling

Wolfgang Gatterbauer

CS3200 Database design (fa22)

https://northeastern-datalab.github.io/cs3200/fa22s3/

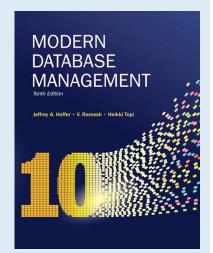
10/31/2022

Class warm-up

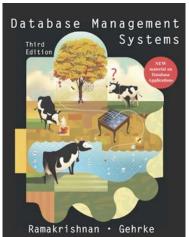
- Last class summary (incl. University example)
- WED: Zoom meeting
 - please use your webcams (it helps me and thus also helps you)
 - slides will be posted WED on Piazza, I am unavailable through SUN
- Project phase 1: see Piazza post, my feedback by TUE
- Exam2: WED next week:
 - SQL part: to be submitted on Gradescope
 - Design: paper & pen: to be handed in afterwards and we bulk-upload
 - normalization may or may not be part of exam (to be decided based on our progress).
 But we continue with normal forms after the exam.
- Thanks for feedback: more in-class text to ERD next time (need to prepare)
- We continue with Database Design, also hands-on



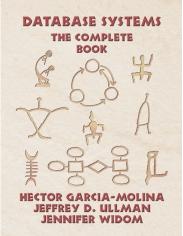
Different sources, different notations



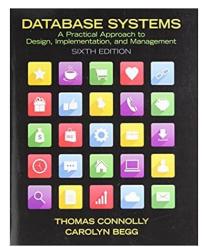




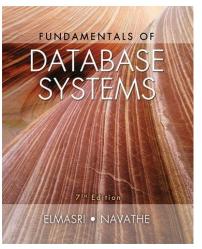
[Cow book'03]

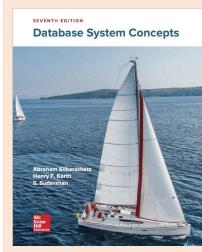


[Stanford book'08]



[Connolly+'15] [Elmasri+'15]





[Silberschatz+'20] SDK arrows

[Hoffer+'10]: Hoffer, Ramesh, Topi. Modern Database Management, 10th ed, 2010.

https://www.pearson.com/us/higher-education/product/Hoffer-Modern-Database-Management-10th-Edition/9780136088394.html

[Cow book'03]: Ramakrishnan, Gehrke, Database Management Systems, 3rd ed, 2003. http://pages.cs.wisc.edu/~dbbook/

[Stanford book'08]: Garcia-Molina, Ullman, Widom. Database Systems: The Complete Book, 2nd ed, 2008. http://infolab.stanford.edu/~ullman/dscb.html

[Connolly+'15]: Connolly, Begg. Database systems: A practical approach to design, implementation, and management, 6th ed, 2015.

https://www.pearson.com/us/higher-education/program/Connolly-Database-Systems-A-Practical-Approach-to-Design-Implementation-and-Management-6th-Edition/PGM116956.html

[Elmasri+'15]: Elmasri, Navathe. Fundamentals of Database Systems, 7th ed, 2015.

https://www.pearson.com/us/higher-education/program/Elmasri-Fundamentals-of-Database-Systems-7th-Edition/PGM189052.html

[Silberschatz+'20]: Silberschatz, Korth, Sudarshan. Database system concepts, 7th ed, 2020. https://www.db-book.com/db7

Wolfgang Gatterbauer. Database design: https://northeastern-datalab.github.io/cs3200/

Notations for binary one-to-many relationships

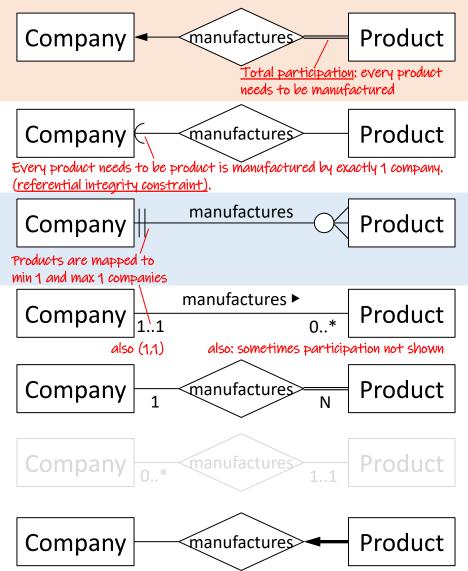
advises

Student

10/24/2022

if you are overwhelmed by the Every student is advised by maximum 1 instructor. different notations, just use An instructor may advise 0, 1 or more students. the one from our textbook SDK [Silberschatz+'20] Instructor advises Student Cardinality constraint: A student is advised by max. 1 instructor (injective) [Stanford book'03] Student Instructor advises also used by Gradiance advises Crow's foot Most often Instructor Student used in practice [Hoffer+'10] Students are advised by min Instructor's advise an look across D and max 1 instructors unrestricted number of students notation advises ▶ **UML** Student Instructor [Connolly+'15] also (D,N) also (0,1) look across for cardinality, [Elmasri+'15] Student Instructor advises same-side for Ν participation Avoid (min-max) !!! same-side Student (min-max) Instructor notation [Elmasri+'15] strongly discouraged since it is the exact opposite of the more commonly used crow's foot look across notation

Every product is manufactured by exactly 1 company. A company may manufacture 0, 1 or more products.

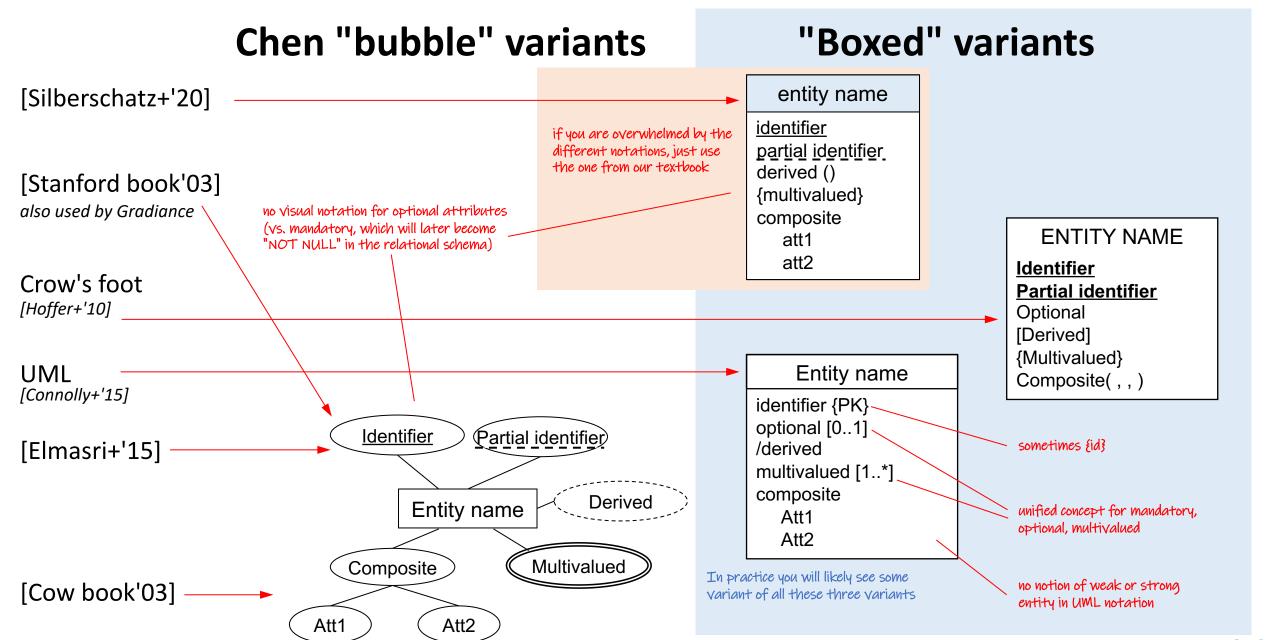


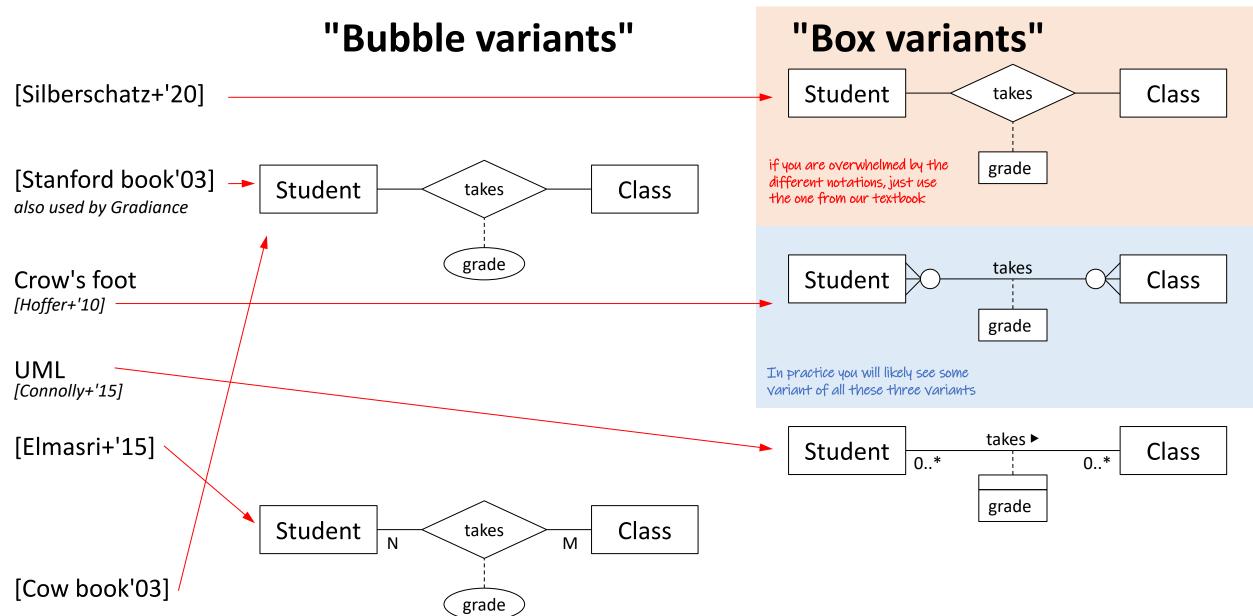
Instructor

[Cow book'03]

if you are overwhelmed by the A course may have 0, 1 or more sections different notations, just use the one from our textbook **Employee** [Silberschatz+'20] depends_on Dependent [Stanford book'03] Employee Dependent depends on also by Gradiance depends on Most often Employee | Crow's foot Dependent used in practice [Hoffer+'10] look across notation depends_on Employee $\frac{1}{1..1}$ **UML** Dependent [Connolly+'15] look across for cardinality, Employee Dependent [Elmasri+'15] depends on same-side for participation Avoid (min-max) !!! same-side (min-max) Employee Dependent notation [Elmasri+'15] strongly discouraged since it is the exact opposite of the more commonly used crow's foot look across notation [Cow book'03] Dependent Employee depends on

Every product is manufactured by exactly 1 company. A company may manufacture 0, 1 or more products. manufactures Product Company Total participation: every product needs to be manufactured (surjective) manufactures Product Company Every product needs to be product is manufactured by exactly 1 company. (referential integrity constraint) manufactures Product Company Products are mapped to min 1 and max 1 companies manufactures > Product Company also (1,1) also: sometimes participation not shown manufacture Product Company Product Company manufactures





Notations for specialization ("ISA relationship")

10/31/2022

(optional=partial | mandatory=total)

[Silberschatz+'20]

Participation (or covering) constraint

if you are overwhelmed by the different notations, just use the one from our textbook

Partial-overlapping optional

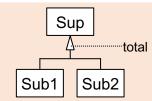
Sup
Super entity
(more generalized,
higher-level)

Sub1 Sub2 Sub entity

Sub entity (more specialized, lower-level) Total-overlapping <mark>mand</mark>atory

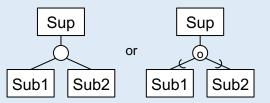
Sup total
Sub1 Sub2

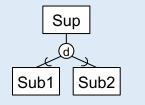
Total-disjoint



[Elmasri+'15]

Crow's foot [Hoffer+'10]



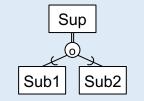


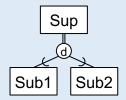
Partial-disjoint

Sub1

Sup

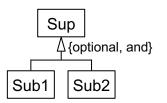
Sub2

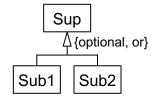


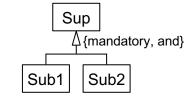


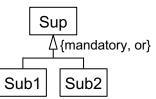
UML [Connolly+'15]

Overlap (or disjoint) constraints (or=disjoint | and=overlapping)





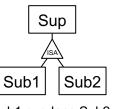


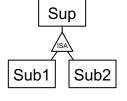


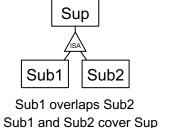
[Stanford book'03]

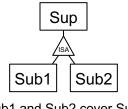
also by Gradiance

[Cow book'03]





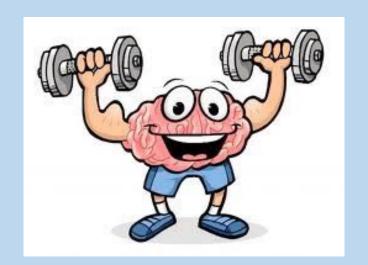




Sub1 overlaps Sub2

verlaps Sub2 Sub1 and Sub2 cover Sup

More Practice



1. Multivalued attributes represented as relationships



COURSE

Course ID

Course_Title

{Prerequisites}

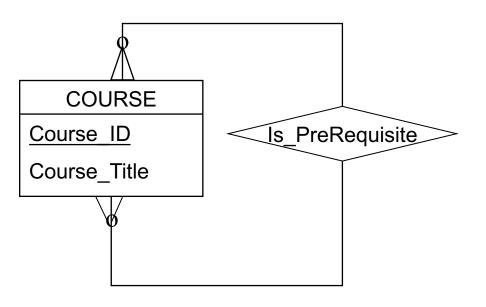




1. Multivalued attributes represented as relationships



COURSE Course ID Course_Title {Prerequisites}



2. Multivalued attributes can be represented as entities



EMPLOYEE

EMPLOYEE_ID
EMPLOYEE_NAME

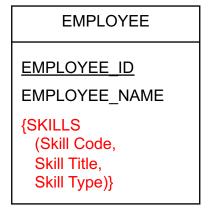
{SKILLS (Skill Code, Skill Title, Skill Type)}

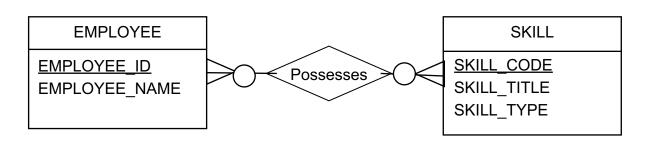
composite

?

2. Multivalued attributes can be represented as entities







composite

3. Attribute vs.



?

EMPLOYEE

EMPLOYEE_ID

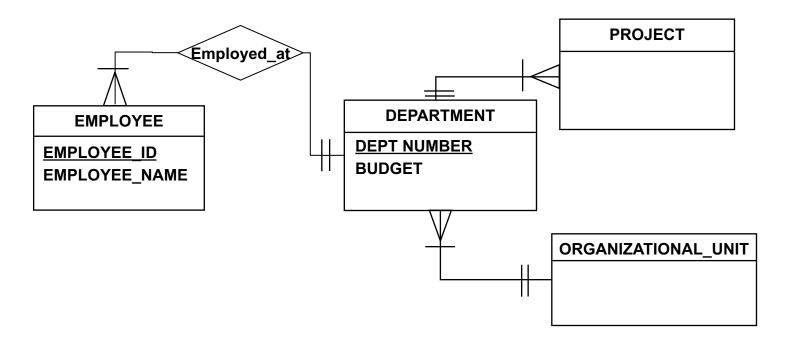
EMPLOYEE_NAME

DEPARTMENT

3. Attribute vs. Multiple Entities



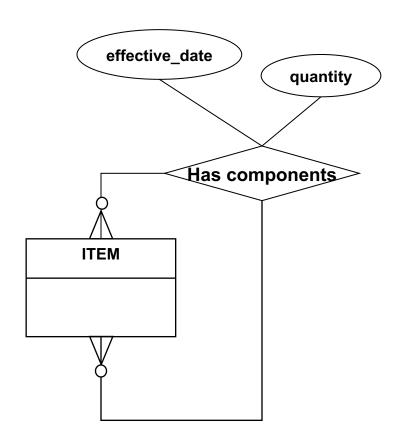




Bill-of-materials (BOM) structure



Relationship

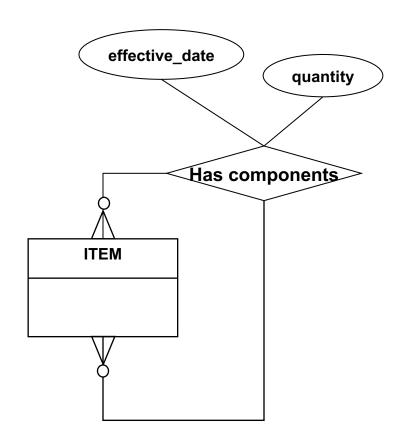




Bill-of-materials (BOM) structure



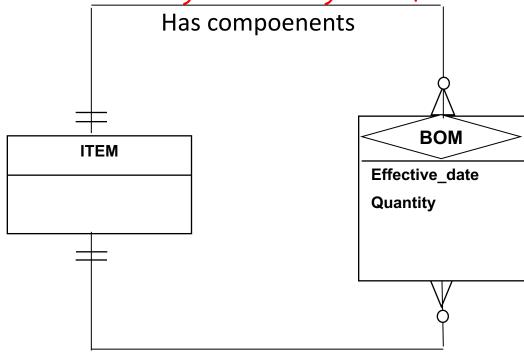
Relationship



Associative entity

possible but I would not recommend; unless want to preserve history:

- weak entity with partial identifier
- or strong with surrogate keys



Used in assemblies

Quick overview: "Relational modeling" From ERDs to Relations

Data modeling and Database Design Process

1. ER Diagram

Conceptual Model:

("technology independent") describe main data items

2. Relational Database Design

Logical Model

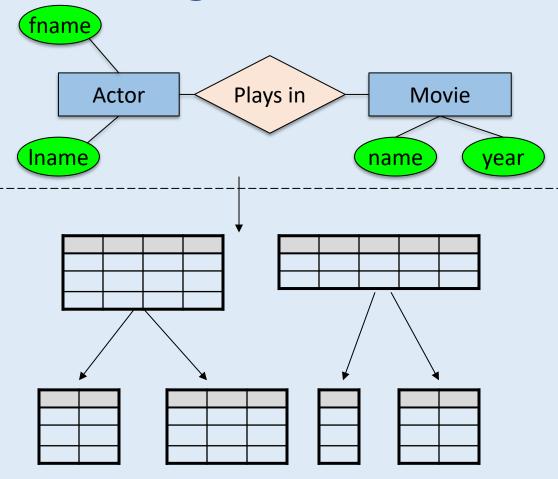
("for relational databases"):

Tables, Constraints

Functional Dependencies

Normalization:

Eliminates anomalies

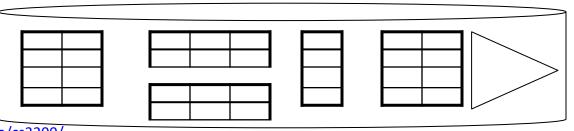


3. Database Implementation

Physical Model

Physical storage details

Result: Physical Schema



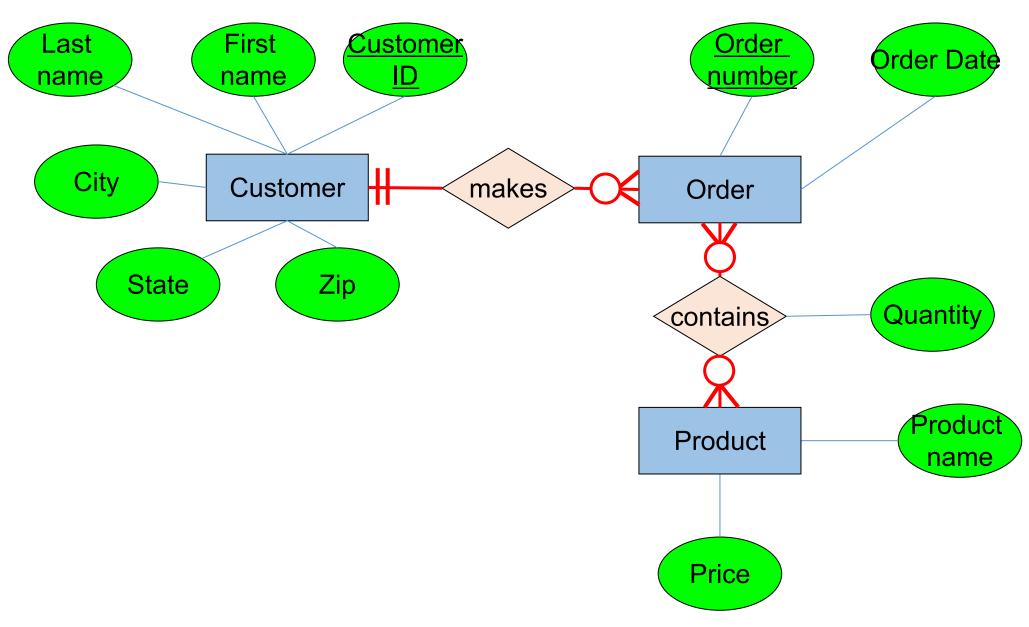
Key concept

- Entity sets become relations, Relationships can become relations (tables in RDBMS)
- Tables are connected with foreign key constraints (~ 1:many relationships!)

A database schema

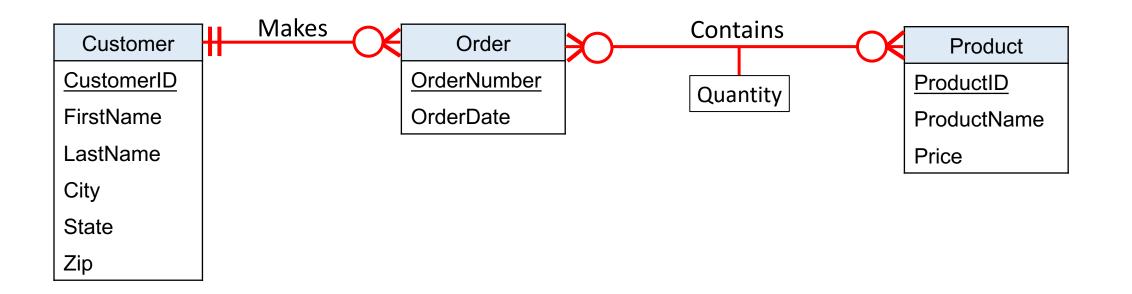
- A map of the tables and fields (attributes) in the database
- This is what is implemented in the database management system
- Part of the "design" process

From ERD to tables (Chen or "bubble" notation)



From ERD to tables (Crow-foot or "block" notation)



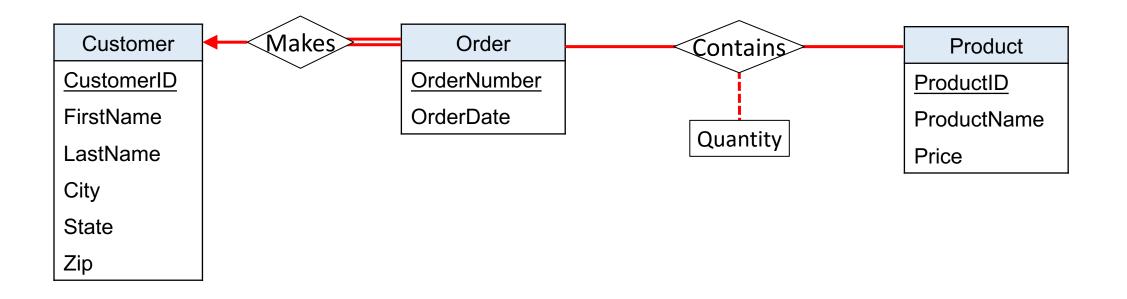


How to translate into our SDK textbook notation?

?

From ERD to tables (our textbok SDK notation)



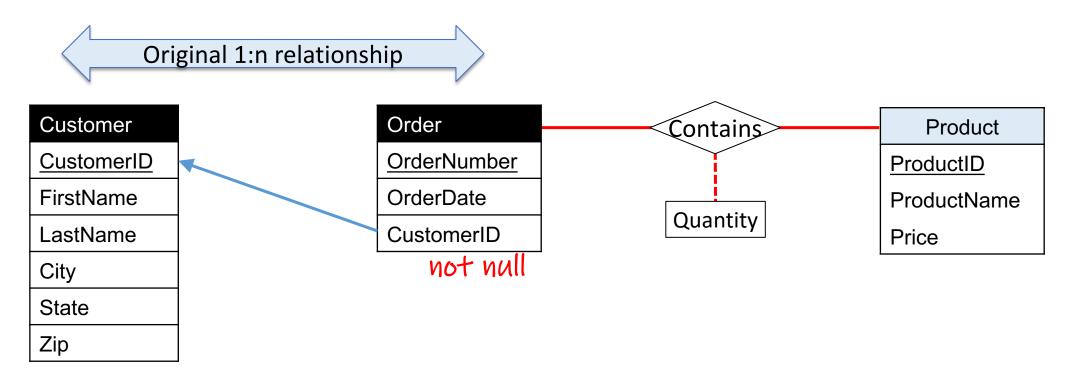


How to translate into tables



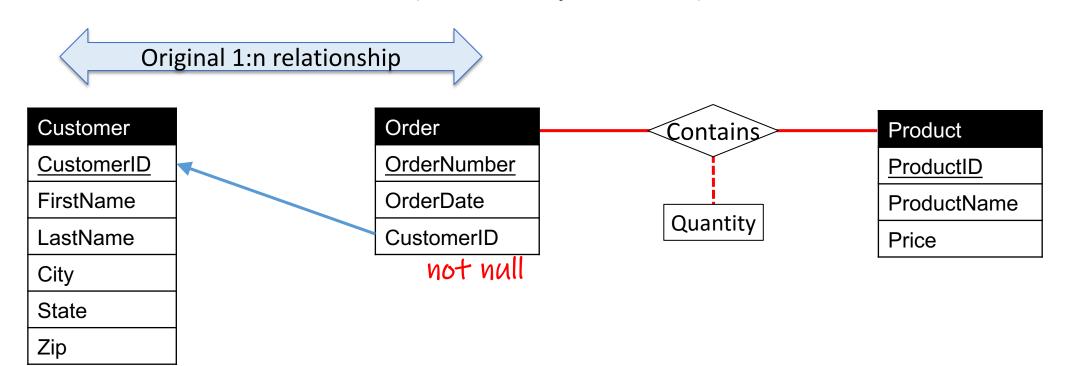
From ERD to tables (incomplete!)





From ERD to tables (incomplete!)



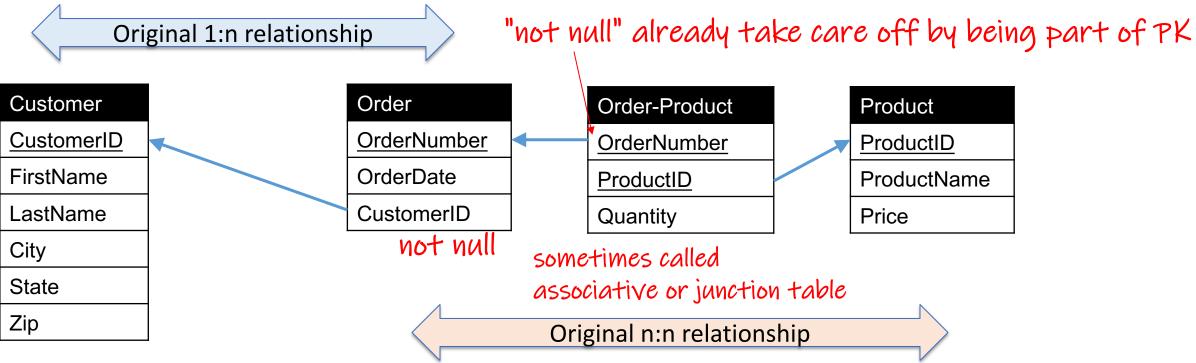


Relational schema

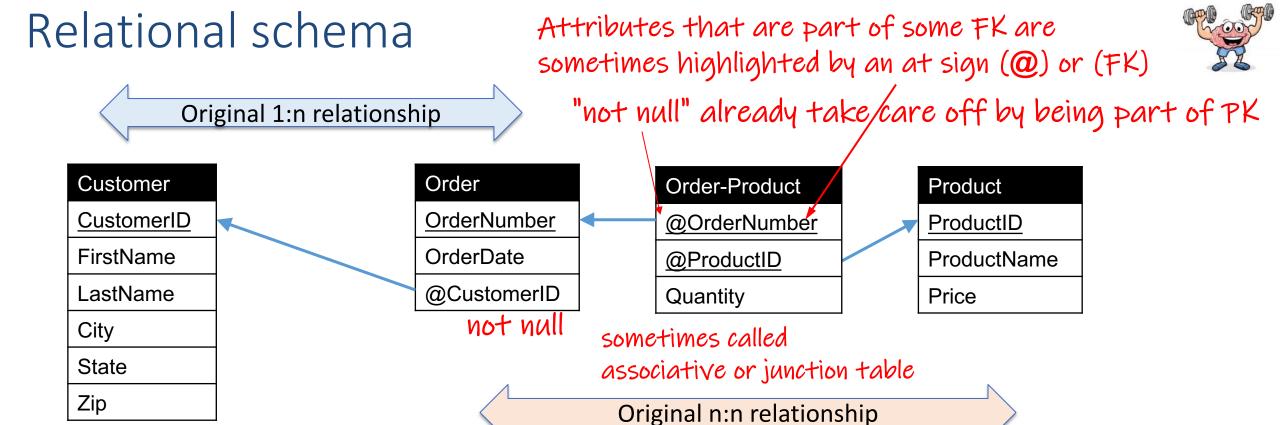




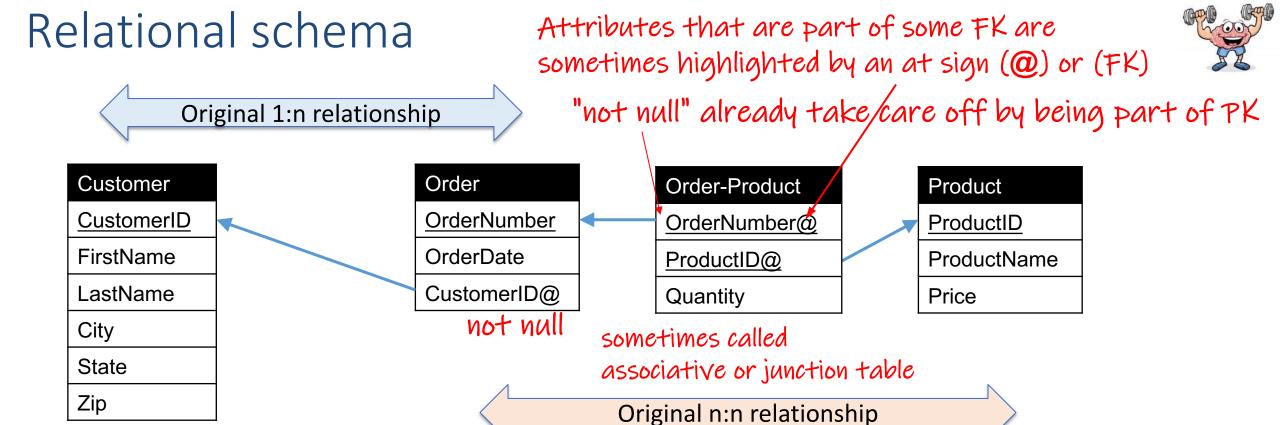




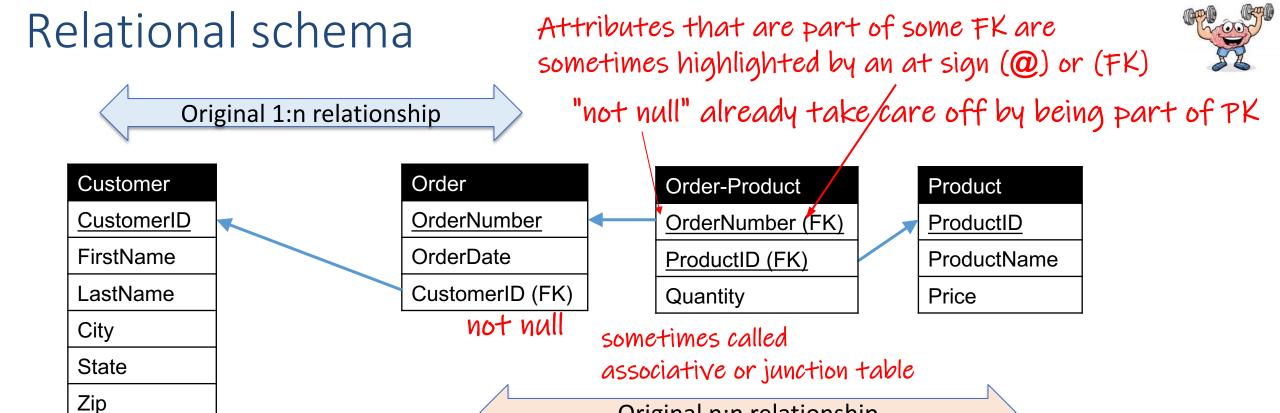
- Order-Product is a decomposed many-to-many relationship
 - Order-Product has a 1:n relationship with Order and Product
 - Now an order can have multiple products, and a product can be associated with multiple orders



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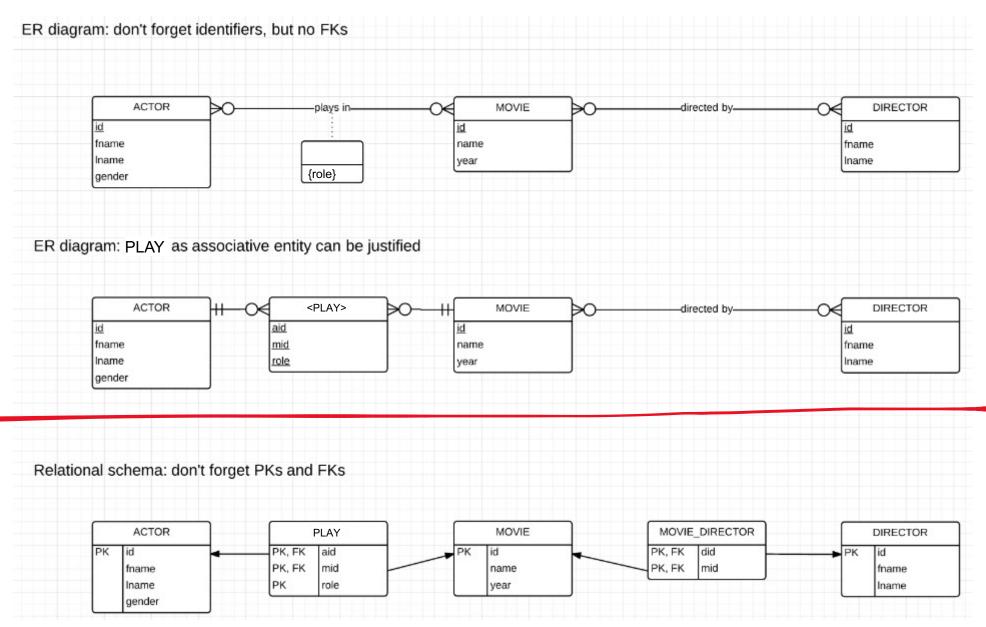
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Original n:n relationship

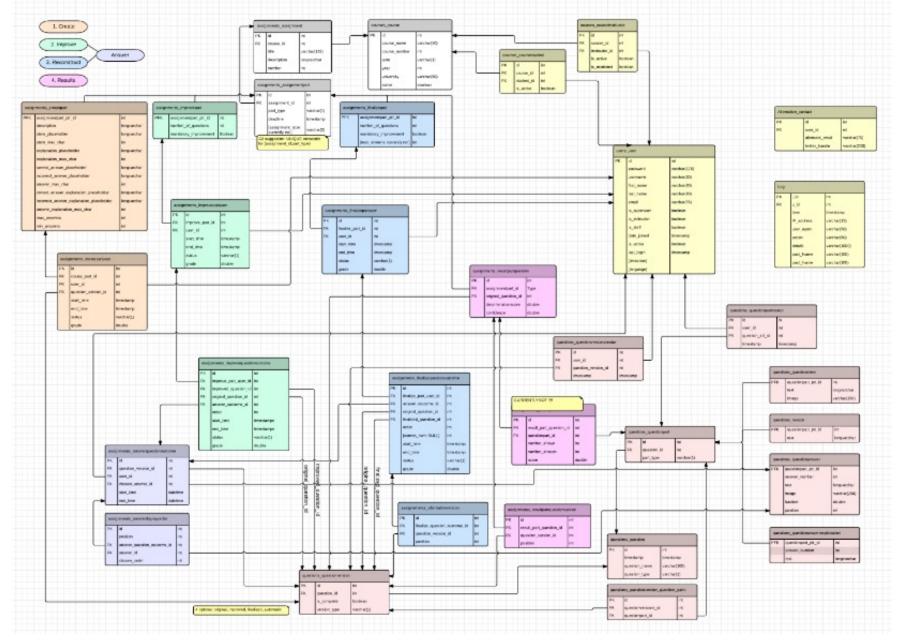
The Rules

- Create a table for every entity
- Create table fields for every entity's attributes
- Implement relationships between the tables
 - 1:many relationships (1:1): primary key field of "1" table put into "many" (other) table as foreign key field
 - many:many relationships:
 - Create new table!
 - 1:many relationships with original table

Play table in our IMDB movie database



Why abstraction? Real relational schemas are huge



Details: "Relational modeling": From ERDs to Relations



Product

ProductID

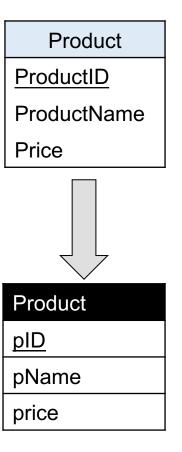
ProductName

Price





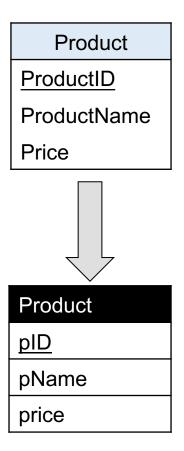
- An entity set becomes a table
 (= relation = multiset of tuples)
 - Each tuple is one entity
 - Each tuple is composed of the entity's attributes, and has the same PK (primary key)





```
CREATE TABLE Product(
pid INT PRIMARY KEY,
pname VARCHAR,
price decimal(8, 2))
```

```
CREATE TABLE Product(
pid INT,
pname VARCHAR,
price decimal(8, 2),
PRIMARY KEY (pid))
```



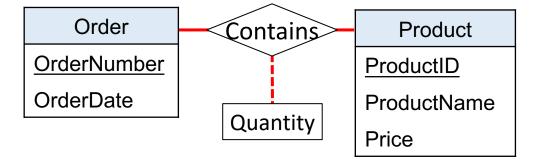
INSERT INTO Product VALUES (1, 'Gizmo', 19.99), (2, 'Supergizmo', 29.99)

Product

| <u>pid</u> | pname | price |
|------------|------------|-------|
| 1 | Gizmo | 19.99 |
| 2 | Supergizmo | 29.99 |

From ER Diagrams to Relational Schema



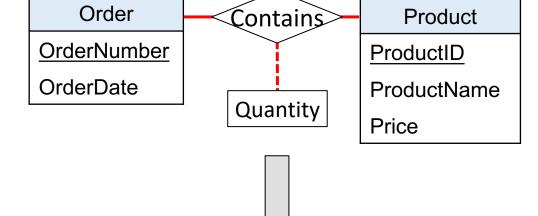




From ER Diagrams to Relational Schema



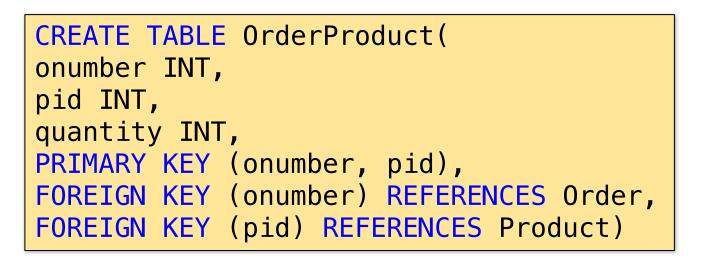
• A many-to-many relationship between entity sets A_1 , ..., A_N also becomes a table (i.e. a multiset of tuples)

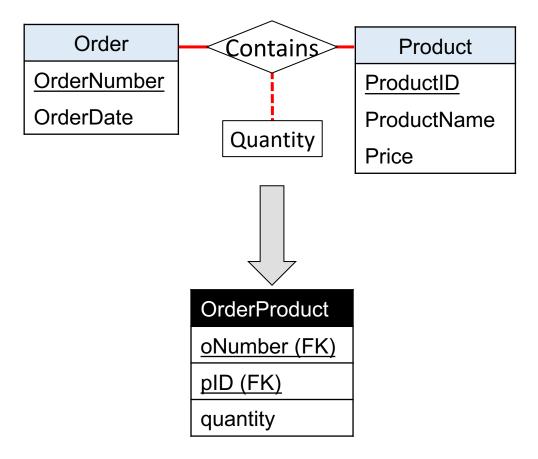


- Each tuple in the table is one relation, i.e. one unique combination of entities $(a_1,...,a_N)$
 - attributes: union of the entity sets' PKs plus the attributes of the relationship
 - FKs (foreign keys): the entities' PKs (primary keys)
 - PK: union of the entity sets' PKs

From ER Diagrams to Relational Schema





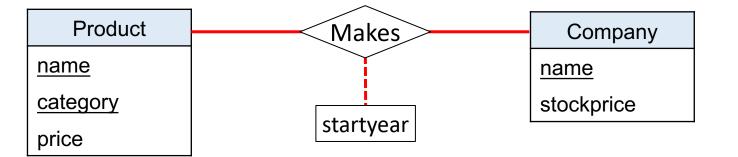


OrderProduct

| INSERT | INT | Ordernumber VALUES |
|--------|-----|--------------------|
| (456, | 1, | 2), |
| (456, | 2, | 3) |

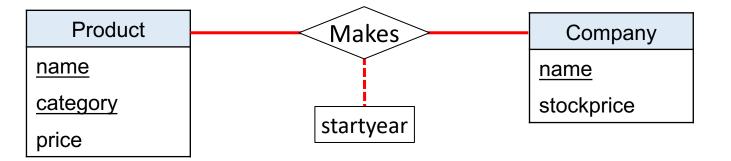
| <u>onumber</u> | <u>pid</u> | quantity |
|----------------|------------|----------|
| 456 | 1 | 2 |
| 456 | 2 | 3 |

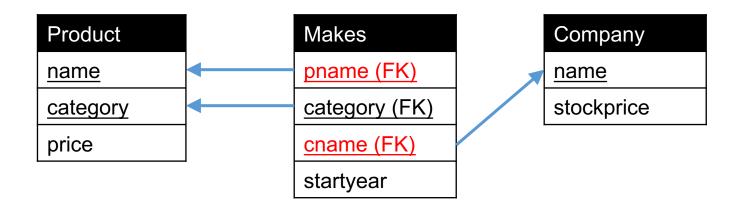






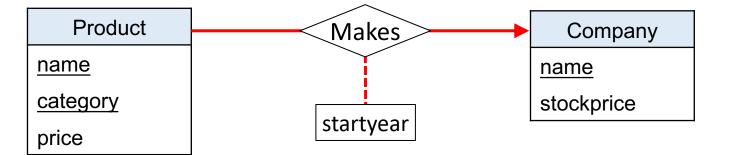






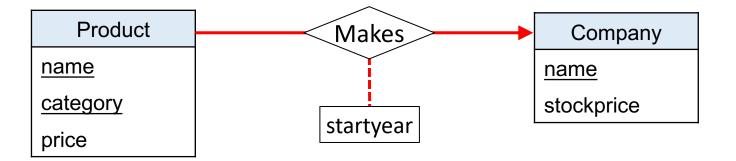
Watch out for attribute name conflicts! Need to rename the FKs.

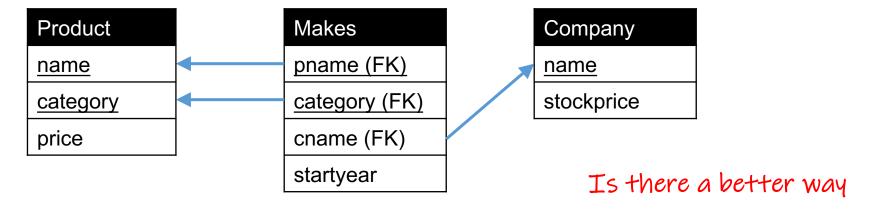






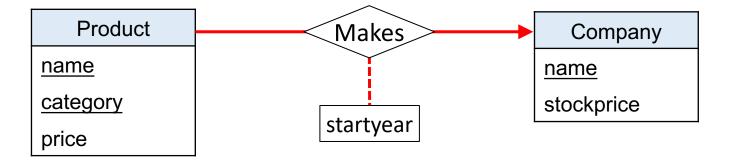


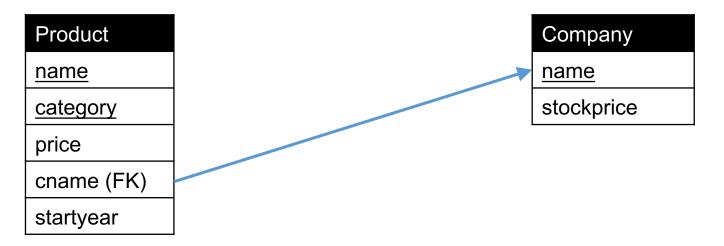




Only keep Product PK as PK



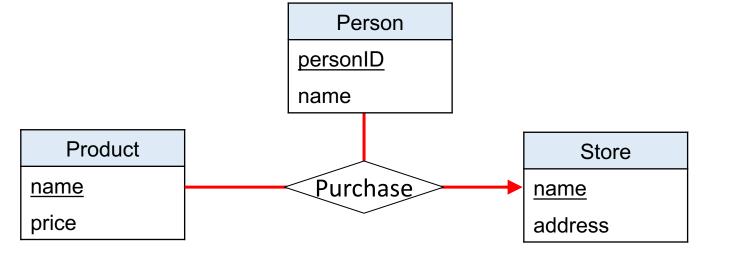




Correct solution: get rid of Makes and add FK and startyear to Product

Multi-way relationships to Relations

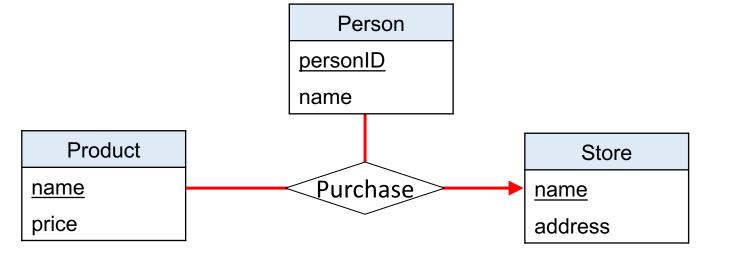






Multi-way relationships to Relations





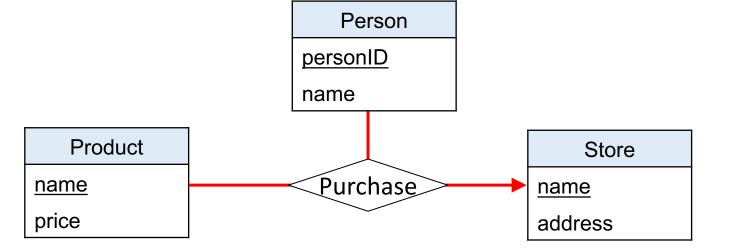
Purchase
pname (FK)
pid (FK)
sname (FK)

What should be the PK?



Multi-way relationships to Relations





Purchase

pname (FK)

pid (FK)

sname (FK)

Notice the composite PK

Exercise (Part 2 & 3): create an ERD & relations



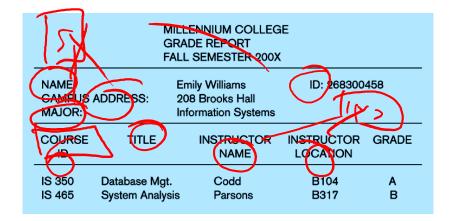
The following grade report below is mailed to students of Millennium College at the end of each semester. Prepare an ERD reflecting the data contained in the grade report (capturing Entities, Attributes, and Relationships). <u>Assume that each course is taught by one instructor and a student can take a class once.</u> Include PKs for each entity type.

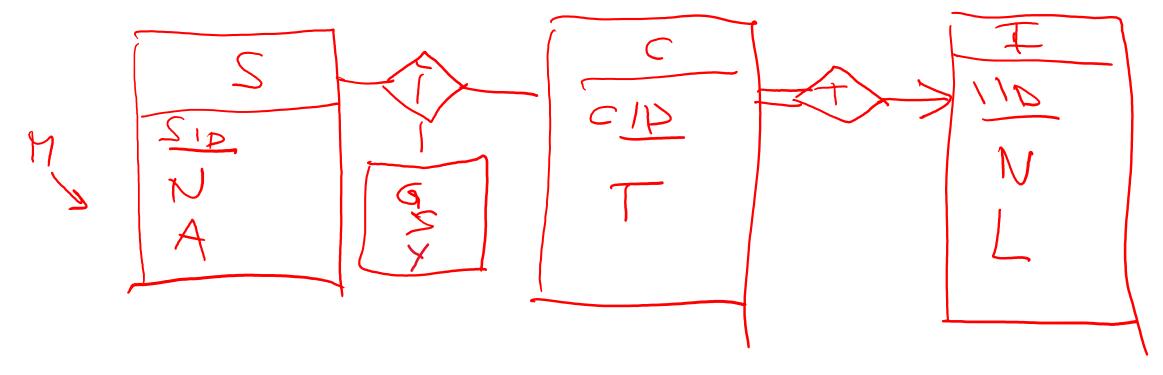
| | GR | LENNIUM COLLEG ADE REPORT LL SEMESTER 200X | _ | |
|-----------------------------|----------------------------------|---|------------------------|--------|
| NAME: CAMPUS / MAJOR: | ADDRESS: 2 | Emily Williams 208 Brooks Hall nformation Systems | ID: 268300 | 458 |
| COURSE | TITLE | INSTRUCTOR NAME | INSTRUCTOR LOCATION | GRADE |
| IS 350 IS 465 | Database Mgt. System Analysis | Codd Parsons | B104 B317 | A B |



Exercise (Part 2): create an ERD





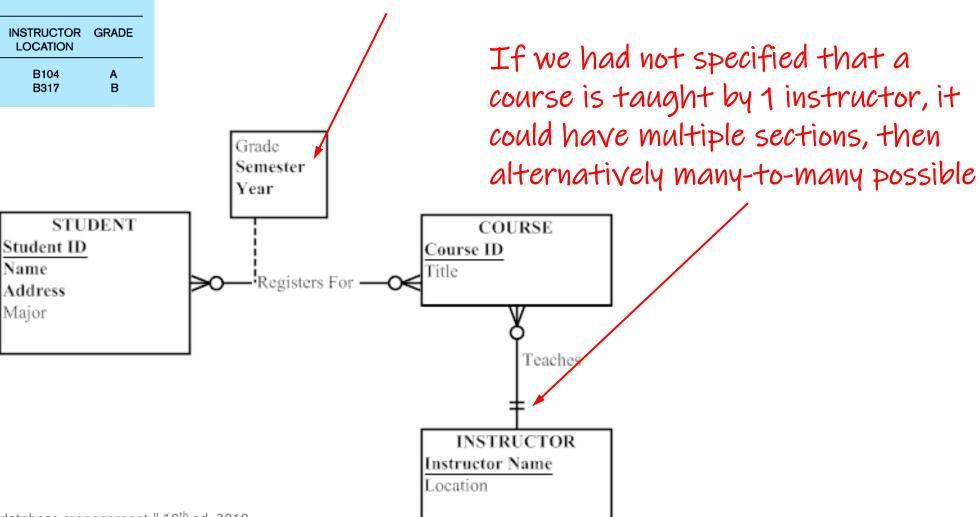


Exercise (Part 2): create an ERD



| MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X | | | | |
|--|----------------------------------|---|------------------------|--------|
| NAME: CAMPUS / MAJOR: | ADDRESS: 2 | mily Williams OB Brooks Hall nformation Systems | ID: 268300- | 458 |
| COURSE | TITLE | INSTRUCTOR NAME | INSTRUCTOR LOCATION | GRADE |
| IS 350 IS 465 | Database Mgt. System Analysis | Codd Parsons | B104 B317 | A B |

Hoffer notation: bold for necessary attributes (not NULL), non-bold for optional ones



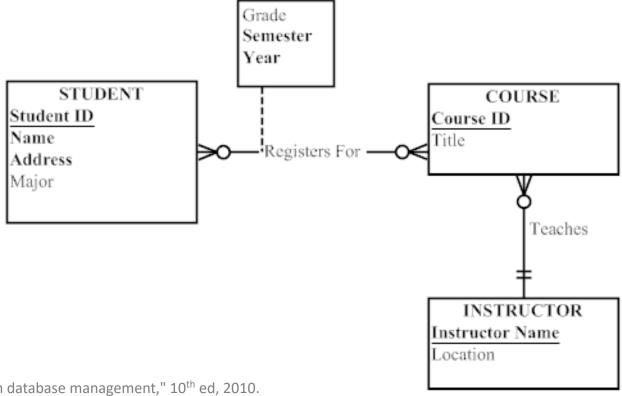
Source: Hoffer, Ramesh, Topi, "Modern database management," 10th ed, 2010. Wolfgang Gatterbauer. Database design: https://northeastern-datalab.github.io/cs3200/

Exercise (Part 3): create relations



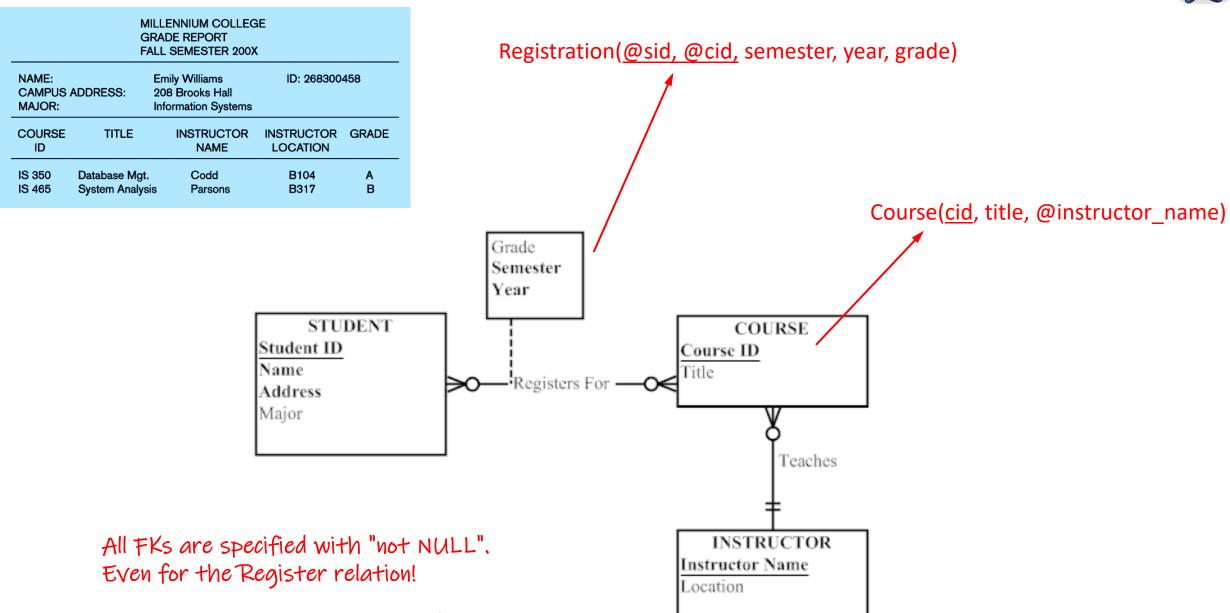
| MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X | | | | |
|--|----------------------------------|--|------------------------|--------|
| NAME: CAMPUS / MAJOR: | ADDRESS: 2 | Emily Williams 208 Brooks Hall Information Systems | ID: 268300 | 458 |
| COURSE | TITLE | INSTRUCTOR NAME | INSTRUCTOR LOCATION | GRADE |
| IS 350 IS 465 | Database Mgt. System Analysis | Codd Parsons | B104 B317 | A B |

Translate into relatoin. Also pay attention to which FKs can be NULL



In-Class Exercise (Part III): relational schema

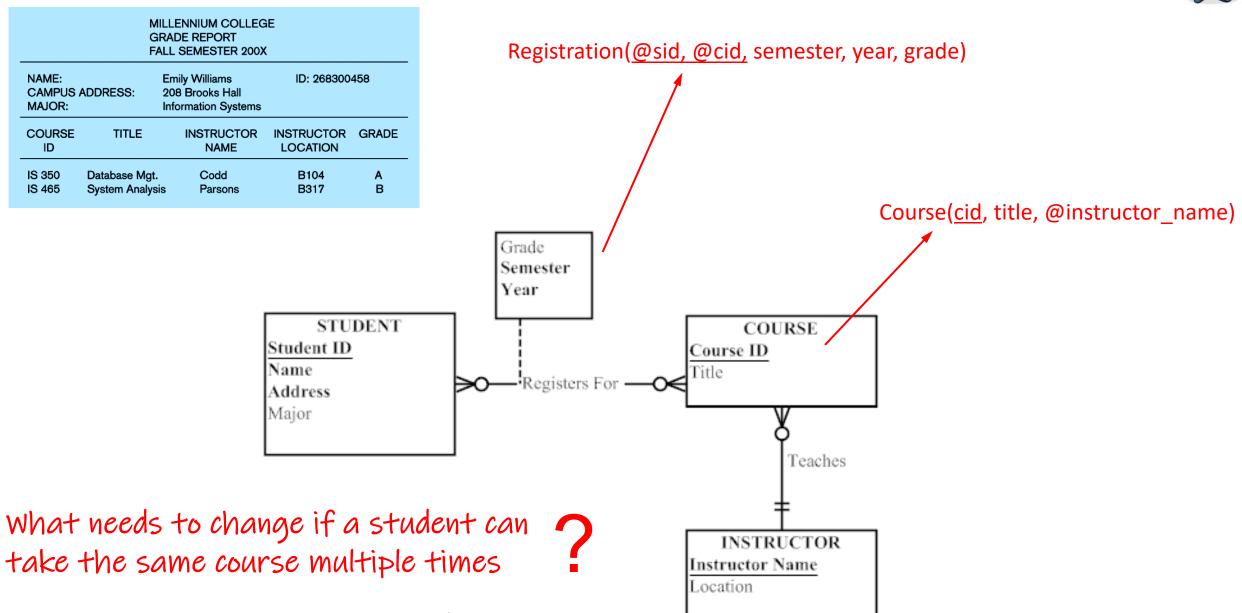




Source: Hoffer, Ramesh, Topi, "Modern database management," 10th ed, 2010. Wolfgang Gatterbauer. Database design: https://northeastern-datalab.github.io/cs3200/

In-Class Exercise (Part III): relational schema



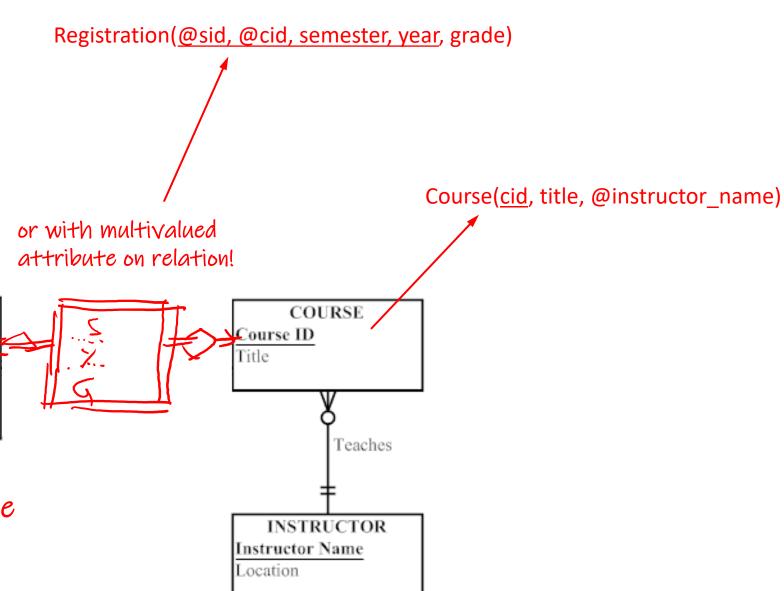


Source: Hoffer, Ramesh, Topi, "Modern database management," 10th ed, 2010. Wolfgang Gatterbauer. Database design: https://northeastern-datalab.github.io/cs3200/

In-Class Exercise (Part III): relational schema



| MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X | | | | |
|--|----------------------------------|--|------------------------|--------|
| NAME: CAMPUS MAJOR: | ADDRESS: | Emily Williams 208 Brooks Hall Information Systems | ID: 268300 | 458 |
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| IS 350 IS 465 | Database Mgt. System Analysis | Codd s Parsons | B104 B317 | A B |



Now a student can take the same course in different semesters

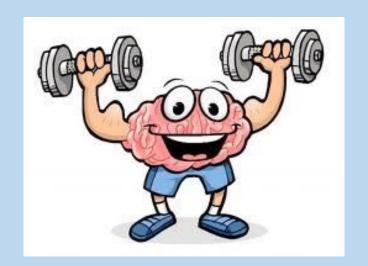
Source: Hoffer, Ramesh, Topi, "Modern database management," 10th ed, 2010. Wolfgang Gatterbauer. Database design: https://northeastern-datalab.github.io/cs3200/

STUDENT

Student ID

Name Address Major

More Practice



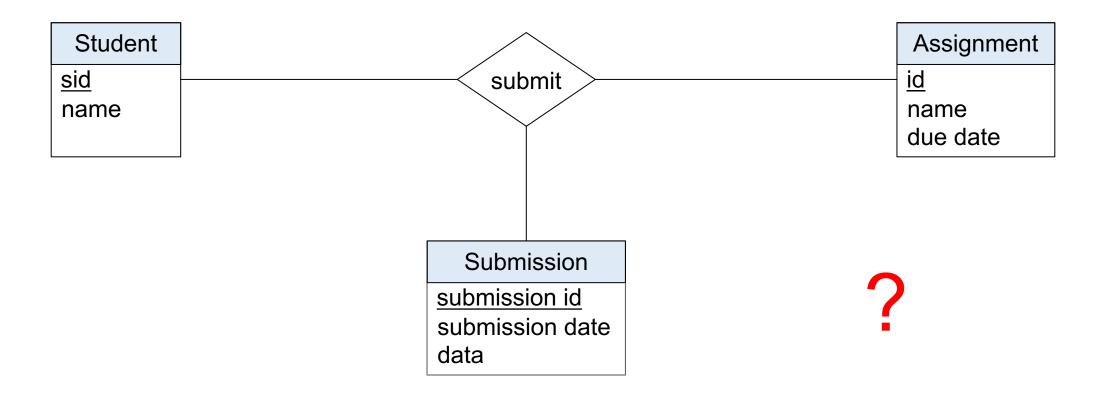


- Students (id, name) submit homework submissions (date, some "data") for assignments (id, name, due date)
- Students can submit multiple versions



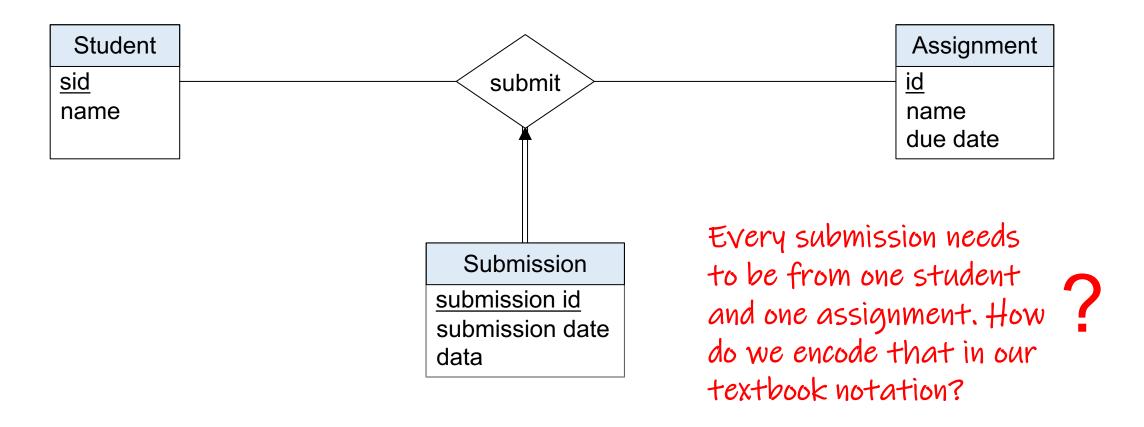


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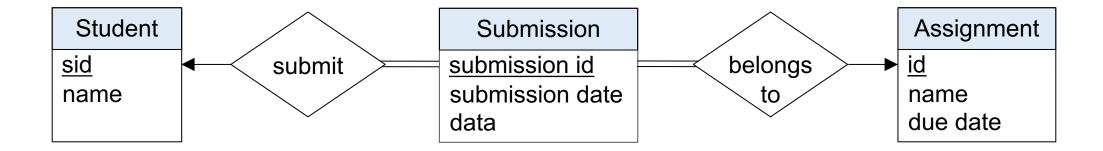


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Assume we want to keep explicity submission versions ?
(1, 2, 3, ...) instead of completely new ids



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- Students can submit multiple versions

