# L12: Relational modeling 2

CS3200 Database design (fa18 s2)

https://northeastern-datalab.github.io/cs3200/

Version 10/18/2018

#### Announcements!

- BB organization cleaned up, no more separate bumping columns
- Continue to come to OHs, read Piazza posts (see Gradiance comment), SAMS book, videos
- Feedback:
  - Great way to share your views, please keep it up!
  - Why the emphasis on in-class participation? ▷ participation teaches an important skill for real life; ▷ there are alternative ways to contribute: e.g., Piazza, "Optional PPTX"; ▷ better than checking class attendance
  - Why do we learn about so many ERD notations? ▷ preparation for real life not just test
- FM exam1:
  - SQL witnesses vs. limits
  - Q1c

# Breaking the "bamboo ceiling" (cp. glass ceiling)



We want <u>you to be successful</u> in real life. Thus, incentives in this class are such to force you to <u>get out of your</u> <u>comfort zone</u> to speak up and to contribute.

"'The loudest duck gets shot' is a Chinese proverb. 'The nail that sticks out gets hammered down' is a Japanese one. Its Western correlative: 'The squeaky wheel gets the grease.'" <u>http://nymag.com/news/features/asian-americans-2011-5/</u> <u>http://en.wikipedia.org/wiki/Bamboo\_ceiling</u> Can you solve witnesses with top-1?

Product (pname, price, cid)



*Q*: Find the most expensive product + its price:

```
SELECTP2.pname, P2.priceFROMProduct P2WHEREP2.price =(SELECT max(P1.price)FROMProduct P1)
```

SELECT pname, priceFROM ProductORDER BY priceLIMIT 1

Can you solve witnesses with top-1?

Product (pname, price, cid)



*Q:* Find the most expensive product + its price:

```
SELECTP2.pname, P2.priceFROMProduct P2WHEREP2.price =(SELECT max(P1.price)FROMProduct P1)
```

SELECT pname, price FROM Product ORDER BY price LIMIT 1

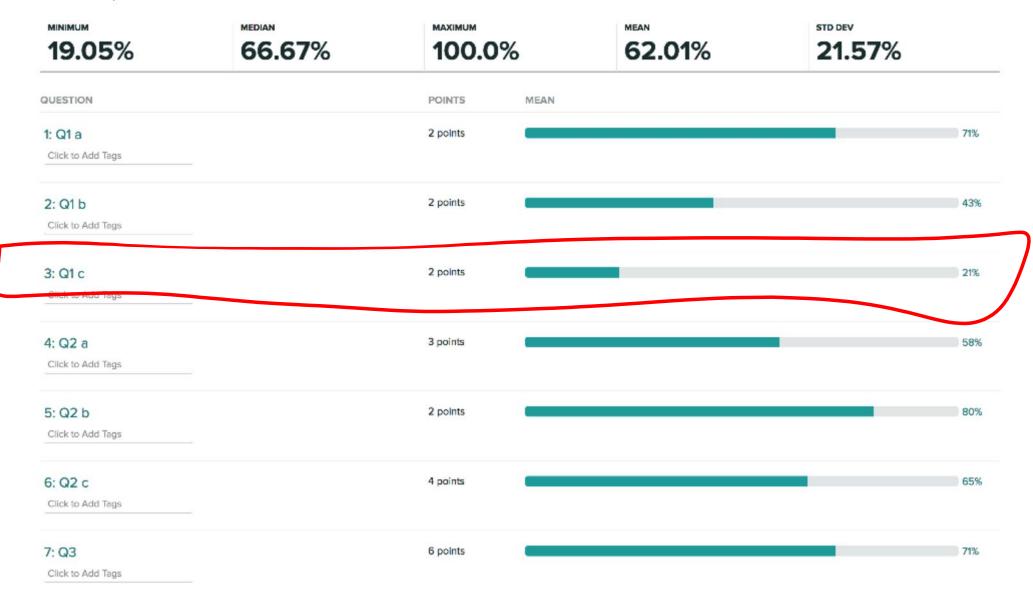
#### Product

Your query needs to return the correct result over \*any\* database

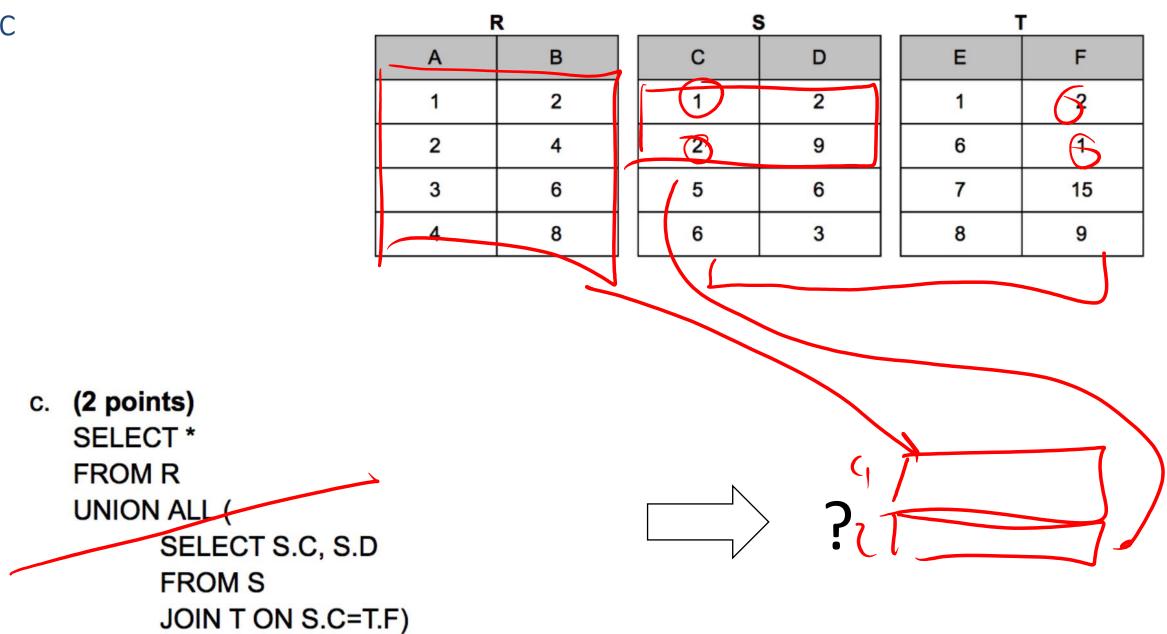
PName	Price	cid
Gizmo	15	1
SuperGizmo	20	1
iTouch1	300	2
iTouch2	300	2

#### Q1-7 on exam 1 (gradescope interface)

Exam1 21.0 points



Q1c

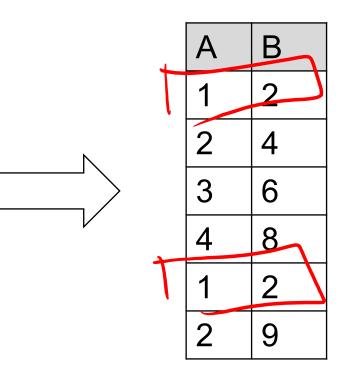


180

Q1c

R		R S		 T	
А	В	С	D	Е	F
1	2	1	2	1	2
2	4	2	9	6	1
3	6	5	6	7	15
4	8	6	3	8	9

c. (2 points) SELECT \* FROM R UNION ALL ( SELECT S.C, S.D FROM S JOIN T ON S.C=T.F)



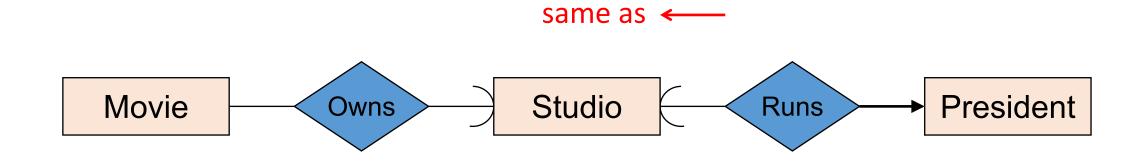
#### Read this ERD to your neighbor





#### Read this ERD to your neighbor





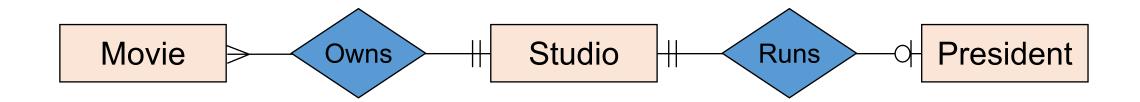
#### Two notations, same meaning



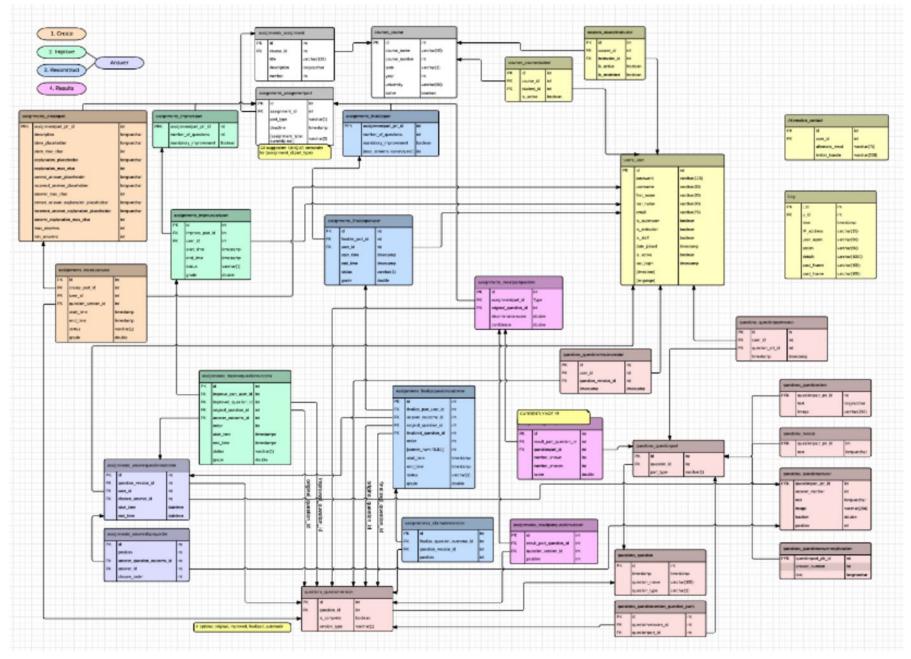


A studio can have at most one president

Each president must run <u>exactly one</u> studio (that exists in the studio entity set)

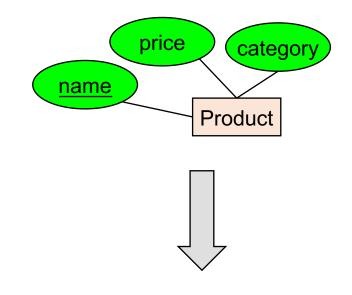


#### Why abstraction? Real relational schemas are huge



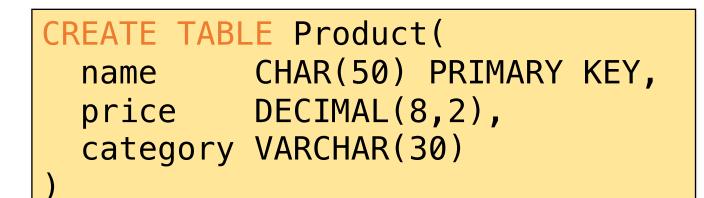
# Back to "Relational modeling": From ERDs to Relations

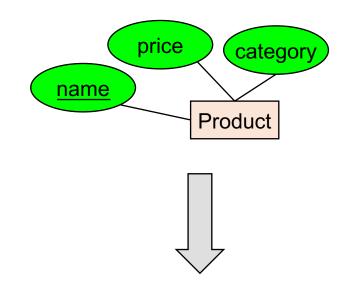
- An entity set becomes a relation (multiset of tuples / table)
  - Each tuple is one entity
  - Each tuple is composed of the entity's attributes, and has the same primary key



Product

name	price	category
Gizmo1	99.99	Camera
Gizmo2	19.99	Edible

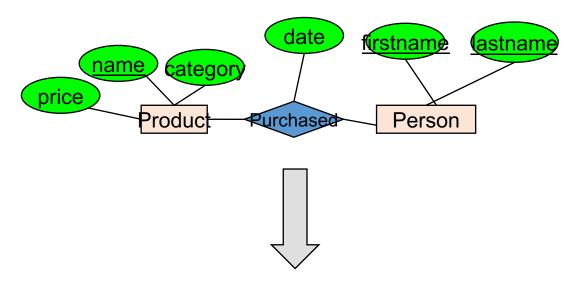




#### Product

name	price	category
Gizmo1	99.99	Camera
Gizmo2	19.99	Edible

- A many-to-many relation <u>between entity sets</u>
   A<sub>1</sub>, ..., A<sub>N</sub> also becomes a multiset of tuples / a table
  - Each row/tuple is one relation, i.e. one unique combination of entities  $(a_1,...,a_N)$
  - Each row/tuple is
    - composed of the union of the entity sets' keys
    - has the entities' primary keys as foreign keys
    - has the union of the entity sets' keys as primary key



#### Purchased

name	firstname	lastname	date
Gizmo1	Bob	Joe	01/01/15
Gizmo2	Joe	Bob	01/03/15
Gizmo1	JoeBob	Smith	01/05/15

```
CREATE TABLE Purchased(

name CHAR(50),

firstname CHAR(50),

lastname CHAR(50),

date DATE,

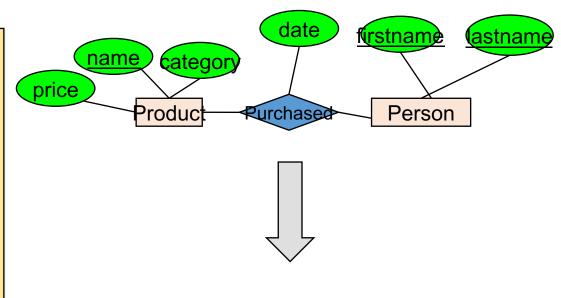
PRIMARY KEY (name, firstname, lastname),

FOREIGN KEY (name)

REFERENCES Product,

FOREIGN KEY (firstname, lastname)

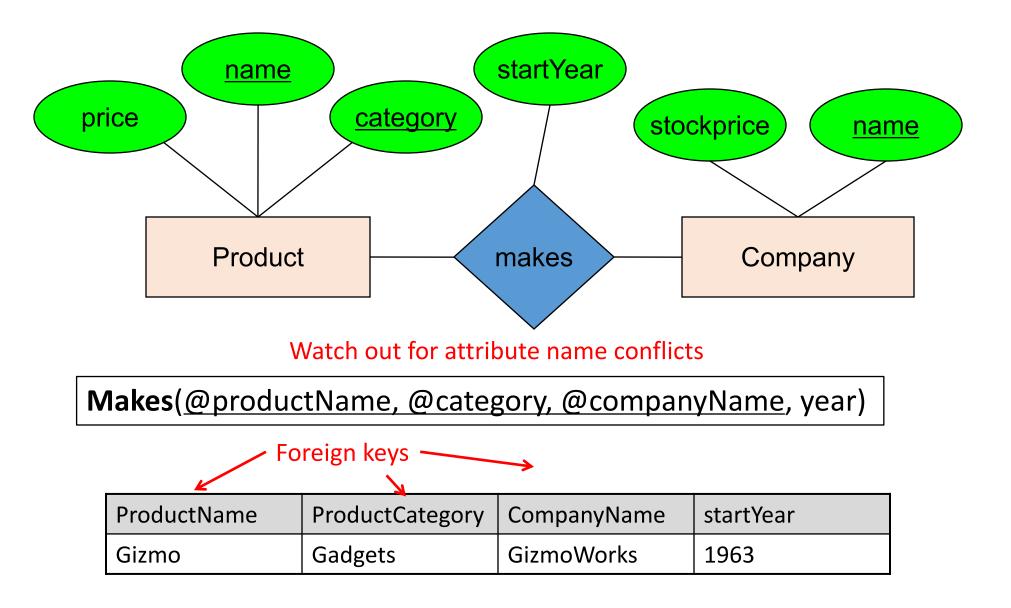
REFERENCES Person
```



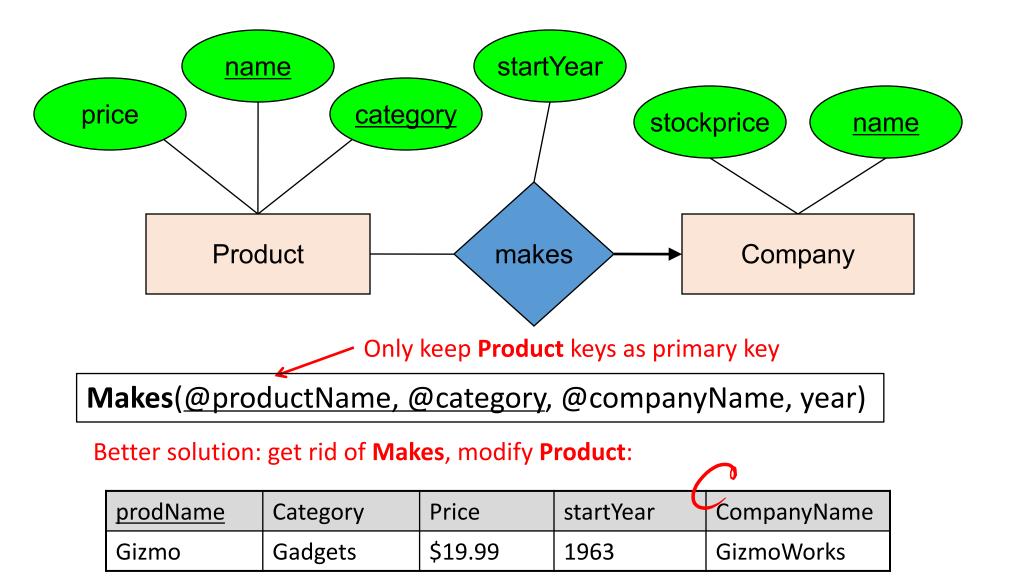
#### Purchased

name	firstname	lastname	date
Gizmo1	Bob	Joe	01/01/15
Gizmo2	Joe	Bob	01/03/15
Gizmo1	JoeBob	Smith	01/05/15

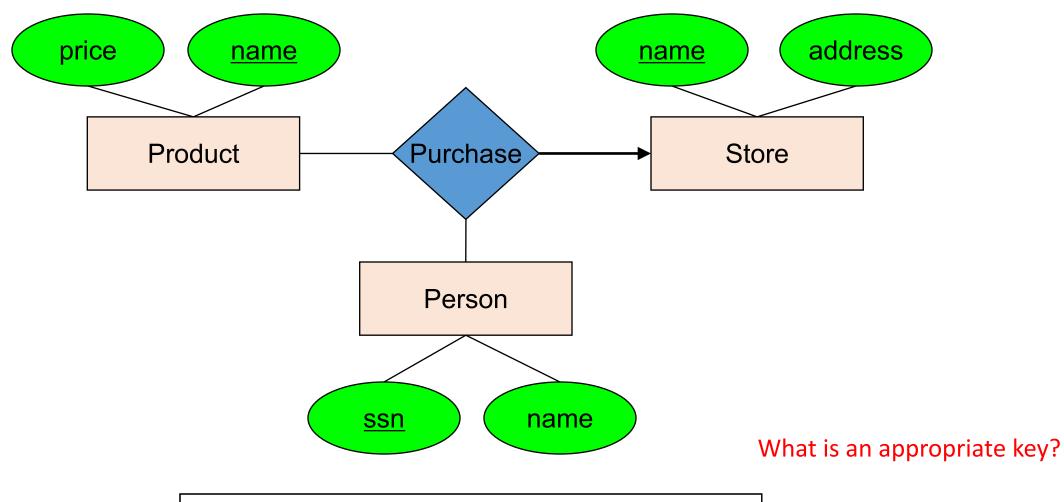
#### Relationships to Relations



### Relationships to Relations (with constraints)

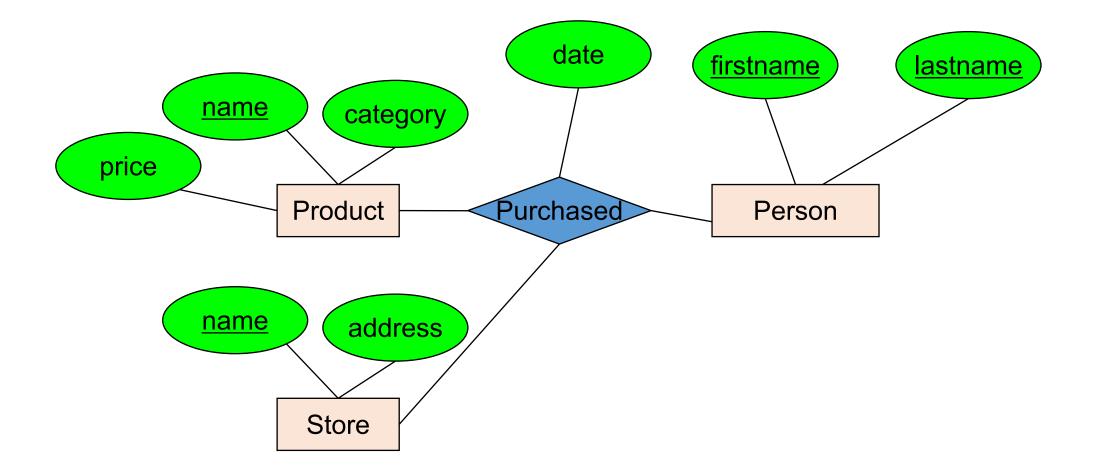


### Multi-way Relationships to Relations



Purchase(prodName, storeName, ssn)

How do we represent this as a relational schema?



#### In-Class Exercise (Part II): create an ERD

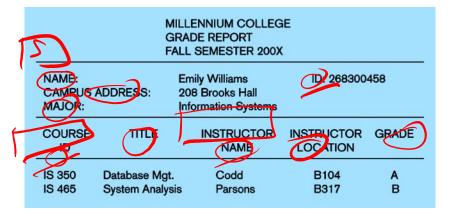


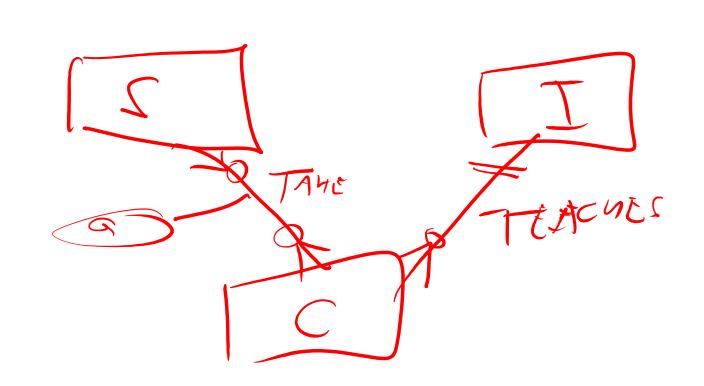
The following grade report below is mailed to students at the end of each semester. Prepare an ERD reflecting the data contained in the grade report (capturing Entities, Attributes, and Relationships). <u>Assume that each course is taught by one</u> <u>instructor</u>. Explain what you chose for the identifier of each entity type

MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X						
NAME: CAMPUS MAJOR:	ADDRESS: 2	Emily Williams 208 Brooks Hall nformation Systems	ID: 268300	458		
COURSE ID	TITLE	INSTRUCTOR NAME	INSTRUCTOR LOCATION	GRADE		
IS 350 IS 465	Database Mgt. System Analysis	Codd Parsons	B104 B317	A B		

#### In-Class Exercise (Part III): create an ERD



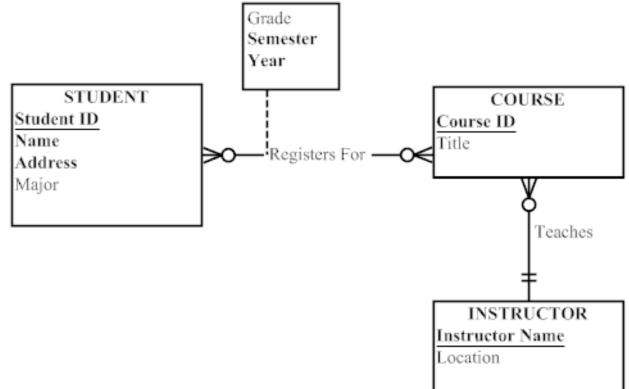




#### In-Class Exercise (Part III): create an ERD

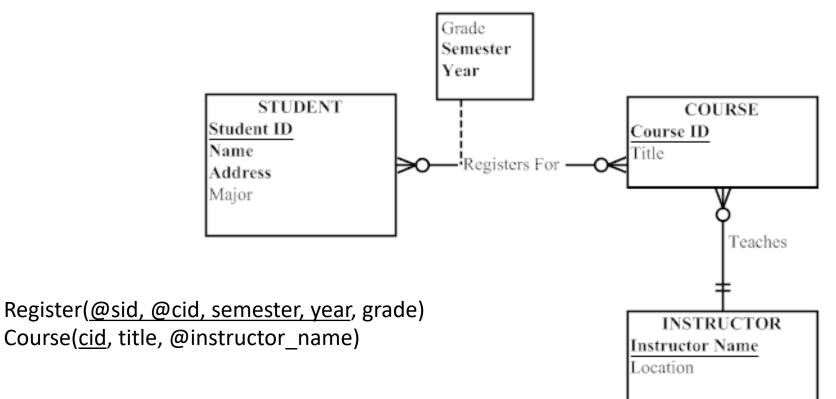
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MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X						
NAME:Emily WilliamsID: 268300458CAMPUS ADDRESS:208 Brooks HallMAJOR:Information Systems						
COURSE ID	TITLE	INSTRUCTOR NAME	INSTRUCTOR LOCATION	GRADE		
IS 350 IS 465	Database Mgt. System Analysis	Codd s Parsons	B104 B317	A B		



#### In-Class Exercise (Part III): relational schema

MILLENNIUM COLLEGE GRADE REPORT FALL SEMESTER 200X					
NAME:     Emily Williams     ID: 268300458       CAMPUS ADDRESS:     208 Brooks Hall       MAJOR:     Information Systems					
COURSE ID	TITLE	INSTRUCTOR NAME	INSTRUCTOR LOCATION	GRADE	
IS 350	Database Mgt.	Codd	B104	А	
IS 465	System Analysis	s Parsons	B317	в	



Relational Modeling in more detail and with some repetition: Entities & Attributes

#### Relations

- A <u>table</u> consists of <u>rows</u> (records), and <u>columns</u> (attributes/fields)
- A <u>relation</u> is a named, two-dimensional table of data
- Six requirements for a table to qualify as a relation:
  - 1. The table must have a unique name.
  - 2. Columns (attributes) in tables must have unique names
  - 3. Every attribute value must be atomic (not multivalued, not composite)
  - 4. Every row must be unique (can't have two rows with exactly the same values for all their columns)
  - 5. The order of the columns must be irrelevant
    - 1. A(<u>id</u>, name) vs. A(name, <u>id</u>)
  - 6. The order of the rows must be irrelevant

### Mapping ER Models To Relations

- Relations (tables) correspond to entity types and to many-to-many relationship types
- Rows correspond to entity instances and to many-to-many relationship instances
- Columns correspond to attributes
- <u>relation (in relational database)</u> ≠ <u>relationship (in E-R model)</u>

EMPLOYEE1						
Emp_ID	Name	Dept_Name	Salary			
100	Margaret Simpson	Marketing	48,000			
140	Allen Beeton	Accounting	52,000			
110	Chris Lucero	Info Systems	43,000			
190	Lorenzo Davis	Finance	55,000			
150	Susan Martin	Marketing	42,000			

#### **Relation Notation**

- Here are two common notations for describing relations:
- Text statements RELATION\_NAME(attributes)
  - CUSTOMER(Customer\_ID, Name, Address, City, State)
  - ORDER(Order\_ID, Order\_Date, Product\_ID)
- Horizontal or Vertical graphical notation:

1	PRODUCT				
	Product_ID	Product_Description	Product_Finish	Standard_Price	Product_Line_ID

### Key Fields

- Keys are special fields used to uniquely identify relations
- Primary keys are unique identifiers of the relation in question
  - Primary keys guarantee that all rows are unique
  - Examples
    - Employee ID numbers
    - Social security numbers
    - E-mail addresses
- Foreign keys are identifiers that refer to other primary keys
  - Useful as references
- Keys can be simple (single field) or <u>composite (multiple fields)</u>
- Keys are often used as indexes to speed up user queries

### Mapping Regular Entities to Relations

- Simple attributes:
  - E-R attributes map directly onto the relation
- Composite attributes:
  - Use only their simple, component attributes

- Multivalued Attribute
  - Becomes a separate relation with a foreign key taken from the superior entity

#### Map Simple ER Attributes Directly Onto Relation



CUSTOMER entity type with simple attributes

CUSTOMER <u>Customer\_ID</u> Customer\_Name Customer\_Address Postal\_Code

#### Map Simple ER Attributes Directly Onto Relation



CUSTOMER entity type with simple attributes

CUSTOMER <u>Customer\_ID</u> Customer\_Name Customer\_Address Postal\_Code

**CUSTOMER** relation

CUSTOMER	l.		
Customer_ID	Customer_Name	Customer_Address	Postal_Code

### Example: Mapping A Composite Attribute

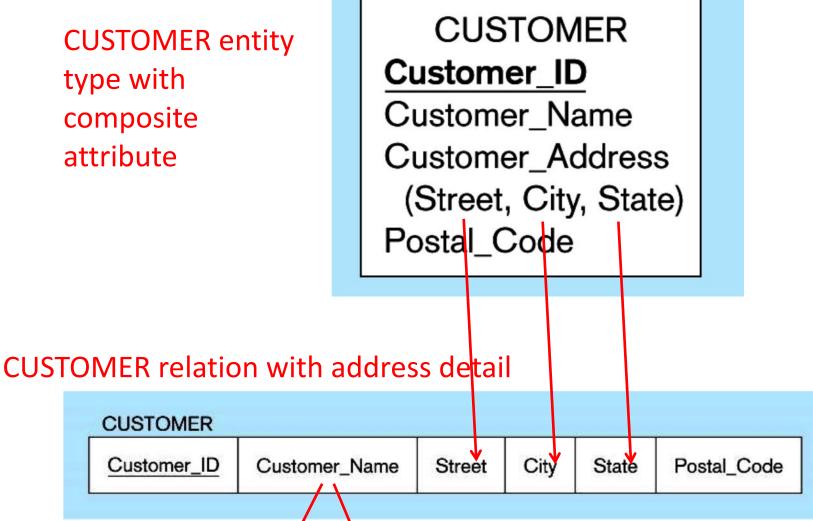


CUSTOMER entity type with composite attribute CUSTOMER <u>Customer\_ID</u> Customer\_Name Customer\_Address (Street, City, State) Postal\_Code

## Example: Mapping A Composite Attribute



**CUSTOMER** entity type with composite attribute



#### Example: Mapping A Multivalued Attribute



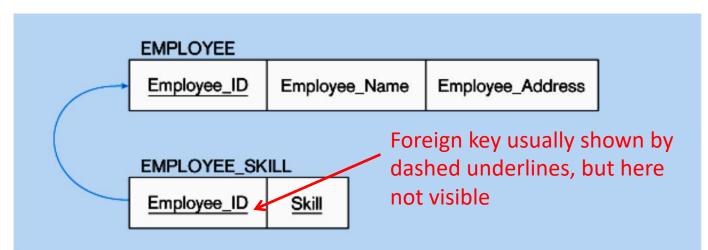
EMPLOYEE <u>Employee\_ID</u> Employee\_Name Employee\_Address {Skill}

# Example: Mapping A Multivalued Attribute



Multivalued Attribute becomes a separate relation with foreign key

EMPLOYEE Employee\_ID Employee\_Name Employee\_Address {Skill}



1-to-many relationship between original entity and new relation

# Relational Modeling: Relationships

# Mapping Binary Relationships

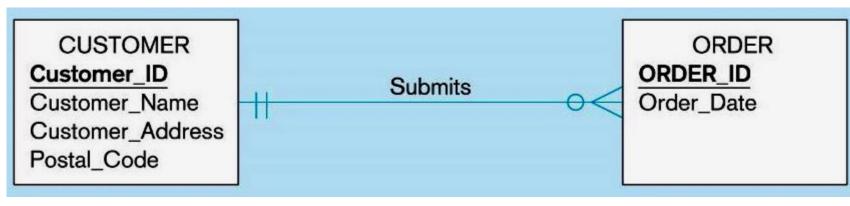
• 1. **One-to-Many**: Primary key on the one side becomes a foreign key on the many side

• 2. Many-to-Many: Create a new relation with the primary keys of the two entities as its primary key

• 3. **One-to-One**: Primary key on the mandatory side becomes a foreign key on the optional side

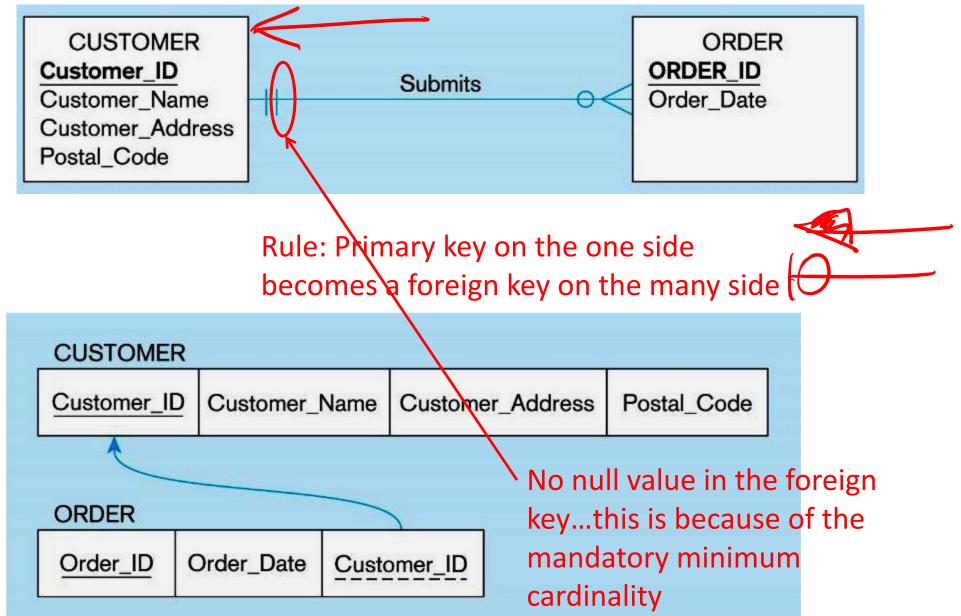
# 1) Mapping a 1:M Relationship





# 1) Mapping a 1:M Relationship







SQLQuery12.sql - jAssignment (1220))*	SQLQuery11.sql - (8OR\Wolfgang (59)	Create the tables
from product;		<pre>create table Company (     CName char(20) PRIMARY KEY,     StockPrice int,     Country char(20) );  create table Product (     PName char(20),     Price decimal(9, 2),     Category char(20),     Manufacturer char(20),     PRIMARY KEY (PName),     FOREIGN KEY (Manufacturer) REFERENCES Company(CName) ); </pre>

	PName	Price	Category	Manufacturer
1	Gizmo	19.99	Gadgets	GizmoWorks
2	MultiTouch	203.99	Household	Hitachi
3	PowerGizmo	29.99	Gadgets	GizmoWorks
4	SingleTouch	149.99	Photography	Canon



SQLQuery12.sql - jAssignment (1220))* SQLQuery11.sql - (8OR\Wolfgang (59)	Create the tables
<pre>insert into product values('hallo', 10, 'Gadgets', NULL);</pre>	<pre>© create table Company (         CName char(20) PRIMARY KEY,         StockPrice int,         Country char(20) );</pre>
	<pre>Create table Product (     PName char(20),     Price decimal(9, 2),     Category char(20),     Manufacturer char(20),     PRIMARY KEY (PName),     FOREIGN KEY (Manufacturer) REFERENCES Company(CName) );</pre>

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Ē	Query12.sql-j select * from product insert into select * from product delete from where manuf	ct; p produc ct; m produc	ct values('	SQLQuery11	Create the tables 
	6 🔹 < Results 📑 Me	essages			PRIMARY KEY (PName), FOREIGN KEY (Manufacturer) REFERENCES Company(CName) )
_	PName	Price	Category	Manufacturer	
1	Gizmo	19.99	Gadgets	GizmoWorks	
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	PName	Price	Category	Manufacturer	
1	Gizmo	19.99	Gadgets	GizmoWorks	
2	hallo	10.00	Gadgets	NULL	
3	MultiTouch	203.99	Household	Hitachi	
4	PowerGizmo	29.99	Gadgets	GizmoWorks	
		149.99	_	Canon	

3

4

2

3

4

5

PowerGizmo

SingleTouch

PName

Gizmo

hallo

MultiTouch

PowerGizmo

SingleTouch

29.99

149.99

Price

19.99

10.00

203.99

29.99

149.99

Gadgets

Category

Gadgets

Gadgets

Gadgets

Household

Photography

Photography

GizmoWorks

Manufacturer

GizmoWorks

GizmoWorks

Canon

NULL

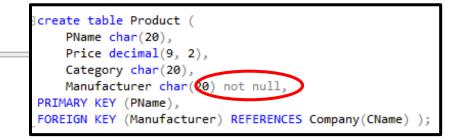
Hitachi

Canon



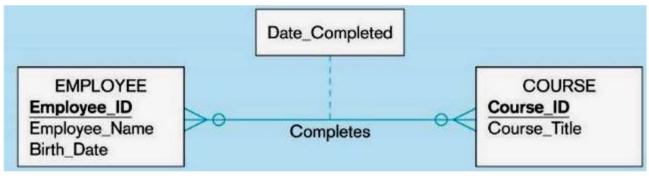
6 100 %	select * from produ delete fro where manu	ict; co produc ict; m produc ifacture:	ct values(' ct	SQLQuery11.	<pre>R\Wolfgang(59) Create the tables create table Company (     CName char(20) PRIMARY KEY,     StockPrice int,     Country char(20) );  create table Product (     PName char(20),     Price decimal(9, 2),     Category char(20),     Manufacturer char(20),     PRIMARY KEY (PName),     FOREIGN KEY (Manufacturer) REFERENCES Company(CName) ); </pre>
	Results 📑 M PName	lessages Price	Category	Manufacturer	
1	_				
	Gizmo	19.99	Gadgets	GizmoWorks	
2	MultiTouch	203.99	Household	Hitachi	What you want

#### What you want:



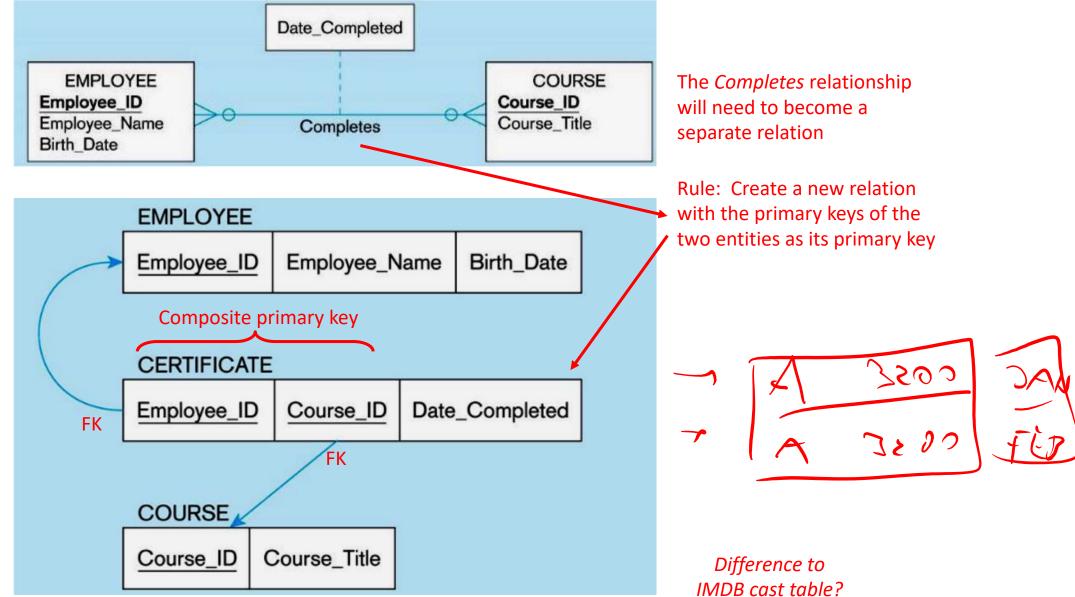
# 2) Mapping An M:N Relationship





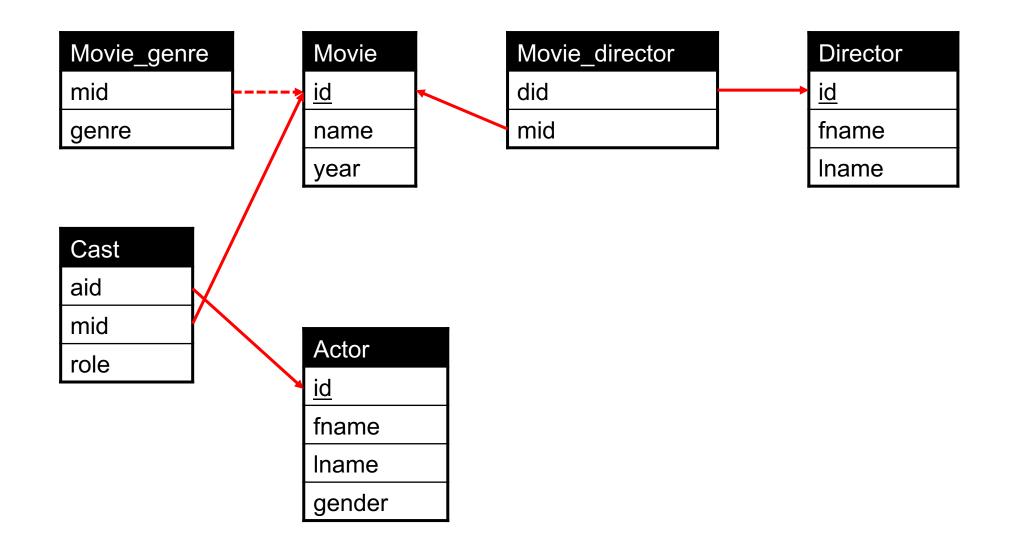
# 2) Mapping An M:N Relationship





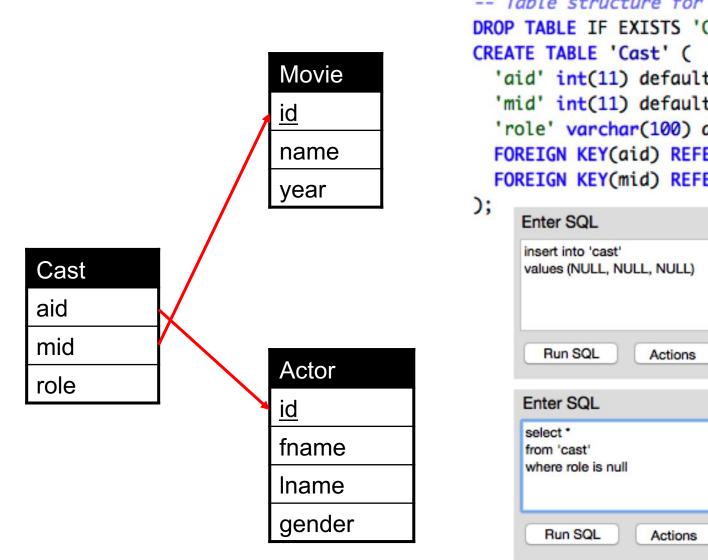
### Small IMDB Movie Database: Schema





#### Small IMDB Movie Database: Schema



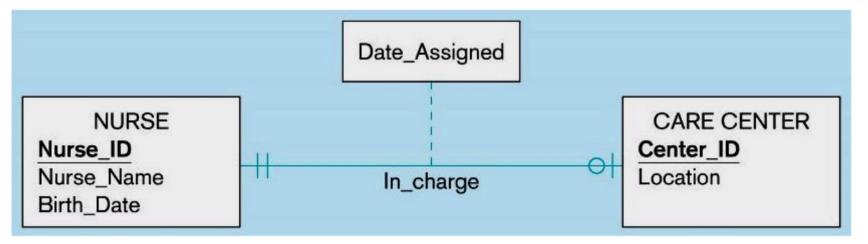


-- Table structure for table 'cast' DROP TABLE IF EXISTS 'Cast'; 'aid' int(11) default NULL, 'mid' int(11) default NULL, 'role' varchar(100) default NULL, FOREIGN KEY(aid) REFERENCES actor(id), FOREIGN KEY(mid) REFERENCES movie(id)

# Last Error: not an error Last Error: not an error aid mid

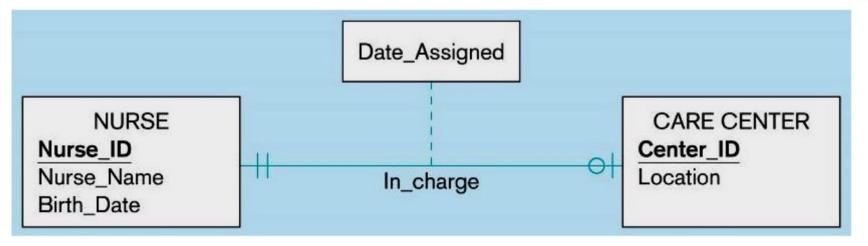
# 3) Mapping A Binary 1:1 Relationship

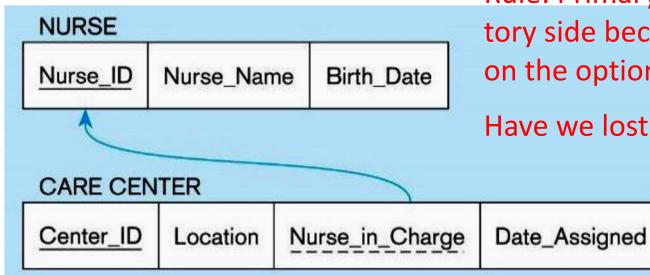




# 3) Mapping A Binary 1:1 Relationship

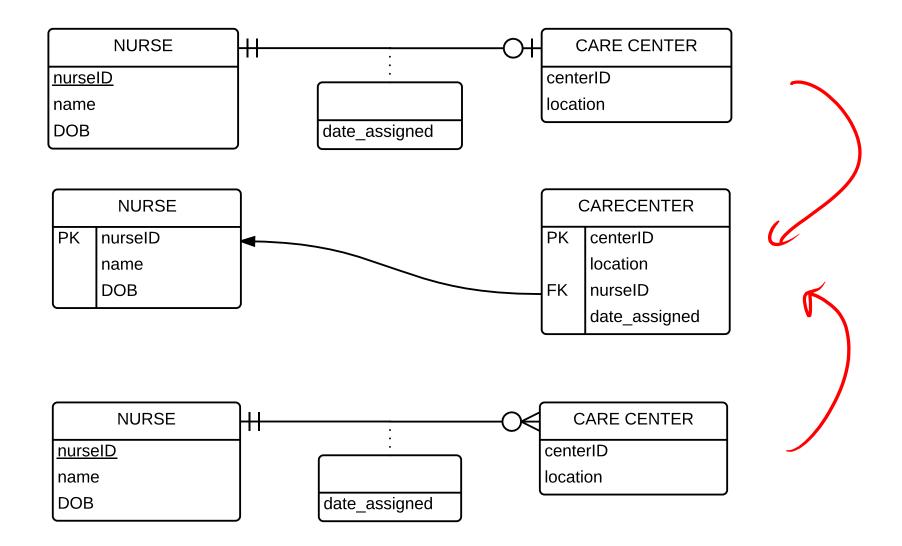




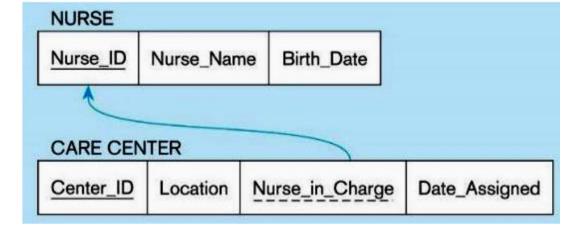


Rule: Primary key on the mandatory side becomes foreign key on the optional side Have we lost some information?

# Transform the ERD into the appropriate schema



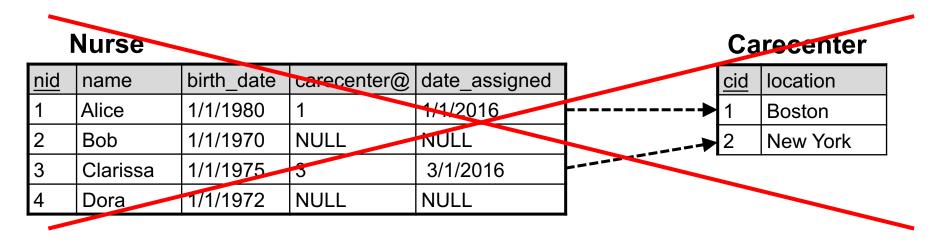
#### Nurses: Instance



#### Nurse

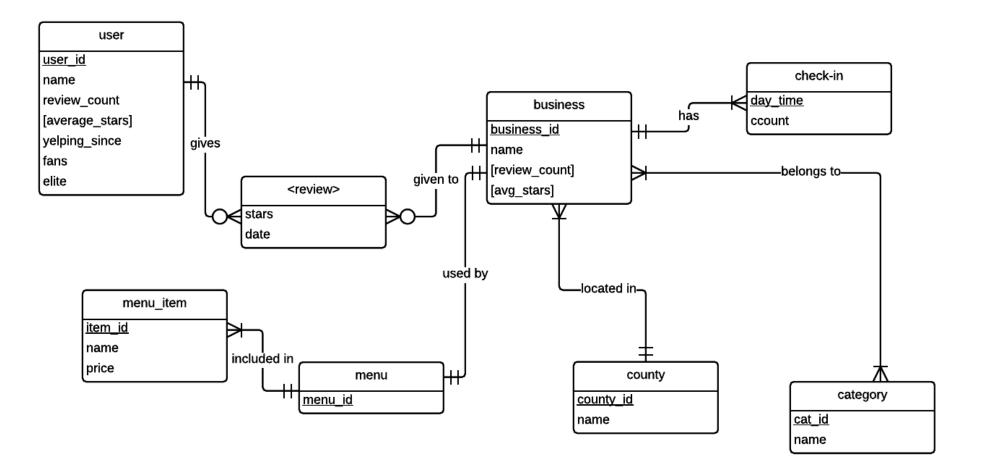
#### Carecenter

nid	name	birth_date		<u>cid</u>	location	nurseid@	date_assigned
1	Alice	1/1/1980	┫	1	Boston	1	1/1/2016
2	Bob	1/1/1970		2	New York	3	3/1/2016
3	Clarissa	1/1/1975				•	
4	Dora	1/1/1972					

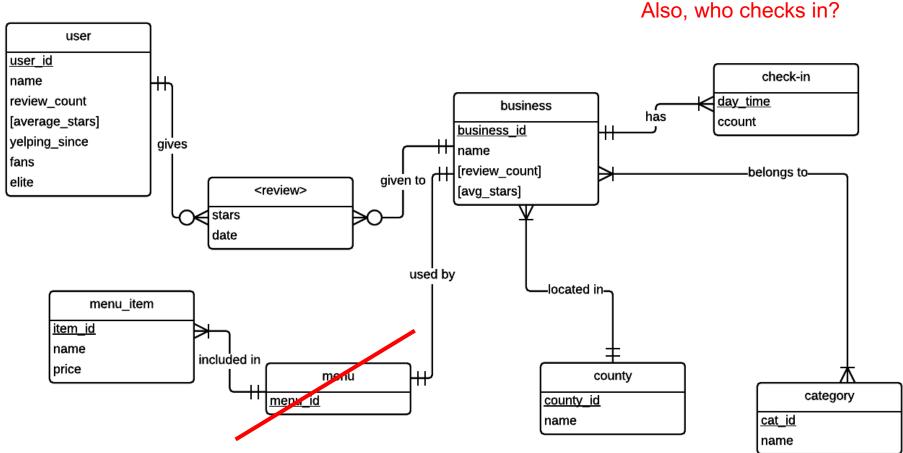


# Yelp: 1:1 relationships

?



### Yelp: 1:1 relationships



No need for separate "menu" given mandatory 1:1 relationship with business. In other words, you will never have a (1,1)-to-(1,1) relationship